

Least Inferable Policies for Markov Decision Processes

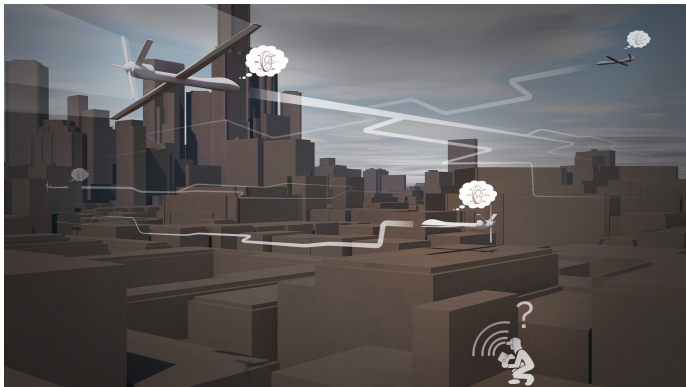
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aUTONOMOUS
SYSTEMS GROUP

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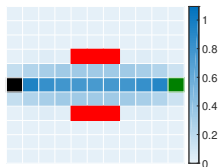
Motivation

An agent that is performing a task in a stochastic environment while being observed by an adversary, should not have an inferable policy.

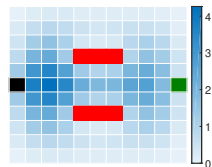


Motivation

Increasing the entropy does not imply non-inferability of the policy.



Least inferable policy



Maximum-entropy policy
(Time limit = 120)

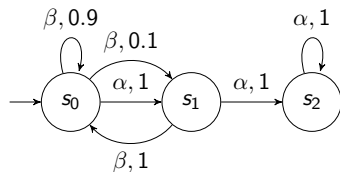
The idea

Synthesize a policy that satisfies some task constraints and limits the ability of the observer to **infer**.

The model

We model the environment with a Markov decision process (MDP) $\mathcal{M} = (S, \mathcal{A}, \mathcal{P}, s_0)$.

- S is a finite set of states,
- \mathcal{A} is a finite set of actions,
- $\mathcal{P} : S \times \mathcal{A} \times S \rightarrow [0, 1]$ is the transition probability function,
- s_0 is the initial state.



A policy is a sequence $\pi = \mu_0 \mu_1 \dots$ where each $\mu_t : S \times \mathcal{A} \rightarrow [0, 1]$ is a function such that $\sum_{a \in \mathcal{A}(s)} \mu_t(s, a) = 1$ for every $s \in S$.

The problem

- The **task constraint** of the agent is to **reach a set S_{reach}** of states with high probability.
- The adversary observes the transitions of the agent at **a set W** of states to **estimate the transition probabilities**.
- The objective of the agent is to minimize **the information leaked** to the adversary on the transition probabilities.
 - What is **leaked information**?

Informativeness of a random variable

The *Fisher information* of a discrete random variable X on θ is

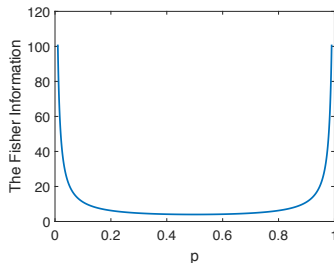
$$I_X(\theta) := \mathbb{E}_X \left[\underbrace{\left(\frac{\partial f(X|\theta)}{\partial \theta} \right)}_{\text{score}}^2 \middle| \theta \right]$$

where $f(X|\theta)$ is the probability mass function.

Example: $Y \sim \text{Bernoulli}(p)$.

$$I_Y(p) = (p(1-p))^{-1}$$

If $p = 0$ or 1 , the inference is easy,
i.e., the estimation error is low.



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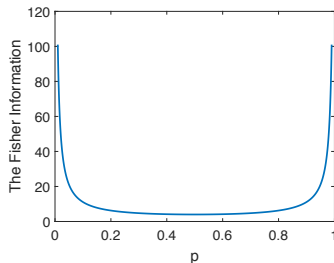
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$X \leftrightarrow$ successor states, $\theta \leftrightarrow$ transition probabilities

Lower bound on the estimation error

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where $f(X|\theta)$ is the probability mass function.

The Cramèr-Rao Bound: Suppose the random variable X is parametrized by θ . The variance of any unbiased estimator $\hat{\theta}$ of θ is lower bounded by the reciprocal of the Fisher information $I_X(\theta)$:

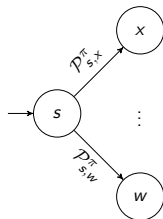
$$\text{Var}(\hat{\theta}) \geq \frac{1}{I_X(\theta)}.$$

Information leaked from a single transition

The *transition information* of a state s is

$$I_s^\pi := \frac{1}{\sum_{q \in S} I_Q(\mathcal{P}_{s,q}^\pi)^{-1}}$$

where Q is the random variable denoting the successor state of state s .



Analogous to Fisher information: Let $\hat{\mathcal{P}}_s$ be an unbiased estimator of the transition probabilities \mathcal{P}_s^π at state s . Then,

$$\text{trace}(\text{Var}(\hat{\mathcal{P}}_s)) \geq \frac{1}{I_s^\pi}.$$

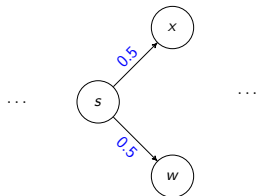
Information leaked from a path

The *total information* of a path $\xi = s_0 s_1 \dots$ is

$$l_{W,\xi}^{\pi} := \sum_{t=0}^{\infty} \mathbb{1}_W(s_t) l_{S_t}^{\pi}.$$

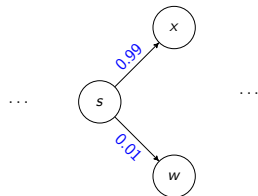
Quantity vs. Informativeness of observations

25 samples from



vs.

1 sample from



Both can be inferred equally well.

The problem

Given

- an MDP $\mathcal{M} = (S, \mathcal{A}, \mathcal{P}, s_0)$,
- a set S_{reach} of states,
- a probability threshold ν_{reach} ,
- the set W of observed states,

compute

$$\min_{\pi}$$

subject to

Expected total information

$$\overbrace{\mathbb{E}_{\xi} [\ell_{W, \xi}^{\pi}]}$$

$$\Pr^{\pi}(\text{Reach}[S_{reach}]) \geq \nu_{reach}.$$

Limiting the policy space

Assumption

The policy π of the agent is stationary, i.e., $\pi = \mu\mu\dots$

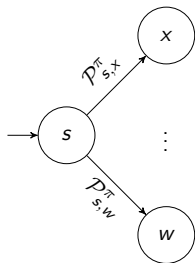
For a stationary policy,

- the expected state residence time at state s is $x_s^\pi = \mathbb{E}[\sum_{t=0}^{\infty} \mathbb{1}_s(s_t)]$,
- the expected state-action residence time at state s and action a is $x_{s,a}^\pi = x_s^\pi \pi_{s,a}$,
- the expected state-state residence time from state s to state q is $y_{s,q}^\pi = \sum_{a \in \mathcal{A}(s)} x_{s,a}^\pi P_{s,a,q}$.

The transition information under a stationary policy π

In terms of transition probabilities:

$$l_s^\pi = \left(\sum_{q \in S} \mathcal{P}_{s,q}^\pi (1 - \mathcal{P}_{s,q}^\pi) \right)^{-1} .$$



In terms of expected residence times:

$$l_s^\pi = \left(\sum_{q \in S} \frac{y_{s,q}^\pi}{x_s^\pi} \left(1 - \frac{y_{s,q}^\pi}{x_s^\pi} \right) \right)^{-1} .$$

A minimum-information admissible policy can be synthesized with a convex optimization problem

$$\min_{x_{s,a}^{\pi}} \sum_{w \in W} x_w^{\pi} l_w^{\pi}$$

$$\text{subject to } l_w^{\pi} = \left(\sum_{q \in S} \frac{y_{w,q}^{\pi}}{x_w^{\pi}} \left(1 - \frac{y_{s,q}^{\pi}}{x_w^{\pi}} \right) \right)^{-1}, \quad \forall w \in W$$

$$x_{s,a}^{\pi} \geq 0,$$

$$\forall s \in S \setminus C, \forall a \in \mathcal{A}(s)$$

$$x_s^{\pi} = \sum_{a \in \mathcal{A}(s)} x_{s,a}^{\pi},$$

$$\forall s \in S \setminus C$$

$$y_{s,q}^{\pi} = \sum_{a \in \mathcal{A}(s)} x_{s,a}^{\pi} \mathcal{P}_{s,a,q},$$

$$\forall s \in S \setminus C, \forall q \in S$$

$$x_s^{\pi} - \sum_{q \in S} y_{q,s}^{\pi} = \mathbb{1}_{s_0}(s),$$

$$\forall s \in S \setminus C$$

$$\sum_{q \in S_{\text{reach}}} \sum_{s \in S \setminus C} y_{s,q}^{\pi} + \mathbb{1}_{s_0}(q) \geq \nu_{\text{reach}}.$$

Expected total information

Flow equations to describe feasible policies

The task constraint

C is the set of the end component states

Lower bound on the estimation error for a state

Let σ_w be the mean-squared error of an (**any**) unbiased estimator for the transition probabilities at state w .

A random path of the agent is the observed data.

Proposition

For an MDP \mathcal{M} and a stationary policy $\pi \in \Pi^{St}(\mathcal{M})$,

$$\sigma_w \geq \frac{(\Pr^\pi(\text{Reach}[w]))^2}{x_w^\pi l_w^\pi}$$

The reachability probability to state w under stationary policy π

The expected leaked information from state w

for every state $w \in W$.

Lower bound on the **total** estimation error

Corollary

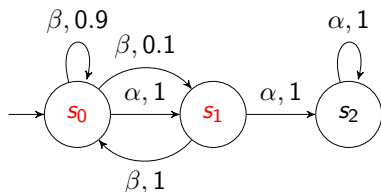
For an MDP \mathcal{M} and a stationary policy $\pi \in \Pi^{St}(\mathcal{M})$,
the total MSE $\sum_{w \in W} \sigma_w$ satisfies

$$\sum_{w \in W} \sigma_w \geq \frac{\min_{w \in W} (\Pr^\pi(\text{Reach}[w]))^2 |W|^2}{\mathbb{E}[\iota_{W,\xi}^\pi]}.$$

The size of set W

The expected
total information

Examples: Estimation error



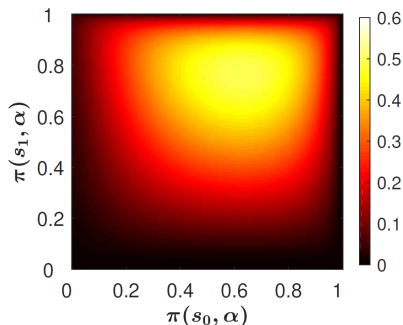
s_0, s_1 : observed

The reachability probability to s_0 and s_1 is 1 under any policy.



Reciprocal of the expected total information \leq

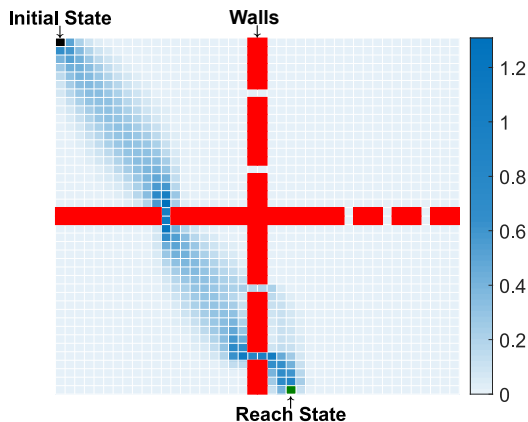
Lower bound on the total mean-squared error.



Total MSE of any unbiased estimator

Examples: Characteristics of the minimum-information admissible policies

Heat map of the expected residence times:

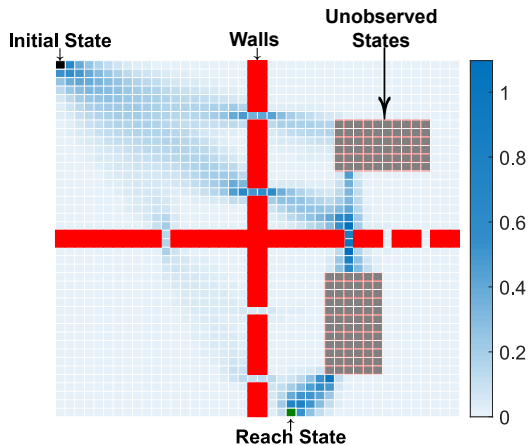


Minimum-information admissible policy yields:

- low number of observations
- less informative observations

Examples: Characteristics of the minimum-information admissible policies

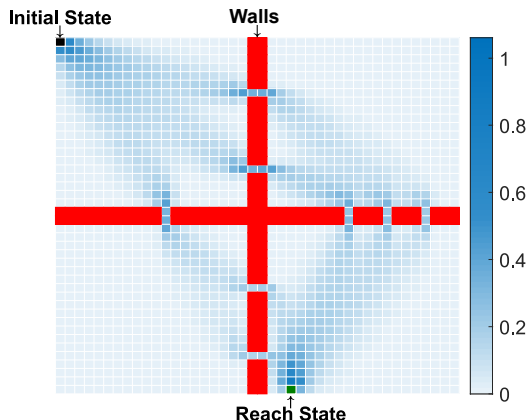
Heat map of the expected residence times:



Minimum-information admissible policy prefers unobserved regions.

Examples: Characteristics of the minimum-information admissible policies with macro-level transition information

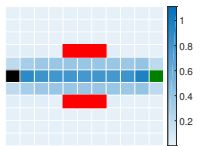
Heat map of the expected residence times:



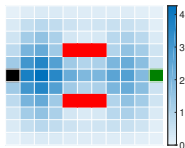
Penalizing the transition information for the gates results in randomization between the gates.

Examples: Comparison of estimation error to maximum-entropy policies

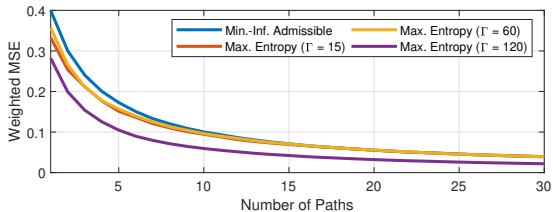
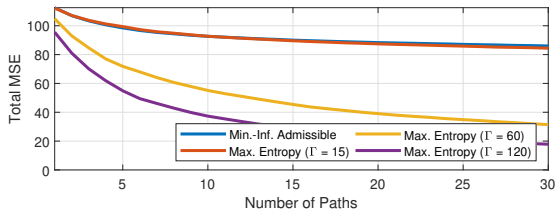
Maximum-entropy policy maximizes the entropy of path distribution given an upper limit Γ on the expected residence times.



Minimum-information admissible policy

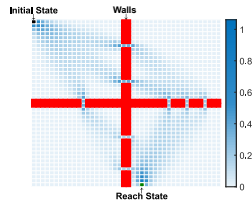
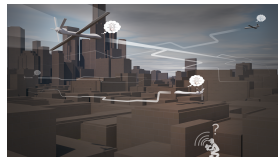


Max. entropy policy (Limit $\Gamma = 120$)



Summary

- Leaked information can be measured with Fisher information
- Computing a minimum-information admissible policy requires to solve a convex optimization problem.
- Estimation error $\propto 1/$ Expected total information



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