## Hyperproperties & Autonomy

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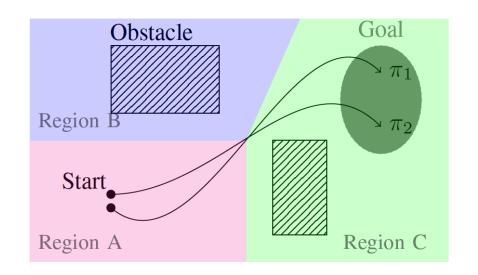
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## **Hyperproperties for Motion Planning**





#### **Motion Planning with Privacy/Opacity**

 $\exists \pi_1 \exists \pi_2. (\pi_1 \text{ and } \pi_1 \text{ are different paths})$   $\land (\pi_1 \text{ and } \pi_1 \text{ give identical observation})$  $\land (\pi_1 \text{ and } \pi_1 \text{ reach goal}).$ 

$$\exists \pi_1 \exists \pi_2. \big( \sec(\pi_1) \neq \sec(\pi_2) \big) \land \big( \operatorname{obs}(\pi_1) = \operatorname{obs}(\pi_2) \big)$$

#### **Optimality of Synthesized Plans**

 $\exists \pi. ((\pi \text{ reaches goal}) \land$ 

 $(\forall \pi'. ((\pi' \text{ reaches goal}) \Rightarrow (\pi \text{ reaches goal})))$ 

$$\exists \pi_1 \forall \pi_2. \ \left( \mathsf{s_0}^{\pi_1} \wedge \mathsf{s_0}^{\pi_2} \right) \wedge \left( \diamondsuit_T (g^{\pi_2} \Rightarrow \diamondsuit_T g^{\pi_1}) \right);$$

$$\exists \pi_1 \forall \pi_2. \ \left( \mathbf{s_0}^{\pi_1} \wedge \mathbf{s_0}^{\pi_2} \right) \wedge \left( \Diamond_T (g^{\pi_1} \Rightarrow \Diamond_T g^{\pi_2}) \right)$$

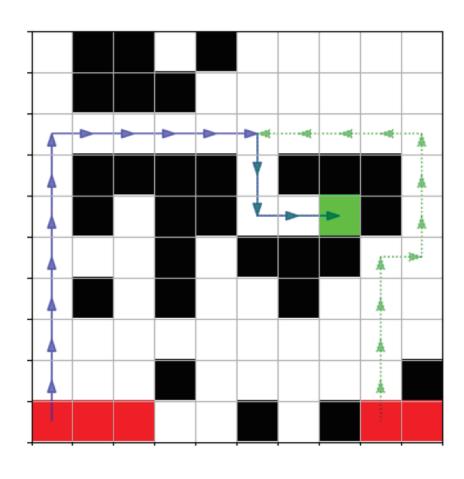
#### **Robustness of Synthesized Plans**

 $\exists \pi \forall \pi'. (\pi \text{ is derived by disturbing } \pi')$  $\land (\pi \text{ and } \pi' \text{ reach goal}).$ 

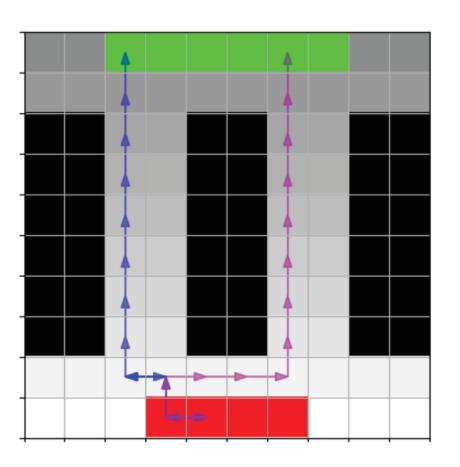
$$\exists \pi_1 \forall \pi_2. \ \operatorname{cls}_{\mathbf{s_0}}(\pi_1, \pi_2) \land \operatorname{cls}_{\mathbf{A}}(\pi_1, \pi_2) \Rightarrow (\varphi^{\pi_1} \land \varphi^{\pi_2})$$

## Symbolic Synthesis from HyperLTL [ICRA'20\*]





**Shortest path** 



**Opacity** 

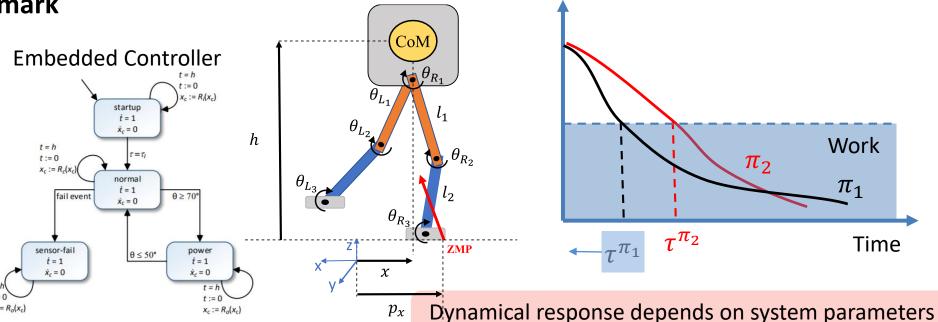
## System Sensitivity to Modeling Errors











#### **Combustion Process**

State	Unit	Description
p	bar	Intake Manifold Pressure
$\lambda_c$	-	A/F Ratio in Cylinder
$\lambda_m$	-	Transfer Function Output
$p_e$	bar	Estimated Manifold Pressure
i	-	Integrator State, PI
$\dot{m}_{af}$	g/s	Inlet Air Mass Flow Rate
$\dot{m}_c$	g/s	Air Flow Rate to Cylinder
$\dot{m}_{\phi}$	g/s	Fuel Mass Aspirated into the Cylinder
$\dot{m}_{\psi}$	g/s	Fuel Mass Injected into Intake Manifold
$\theta_{in}$	degrees	Throttle Angle Input
$\theta$	degrees	Delay-Filtered Throttle Angle
$egin{array}{c}  heta \ \hat{ heta} \end{array}$	-	O/P of Throttle Polynomial
$F_c$	g/s	Command fuel
$\omega$	rad/sec	Engine Speed
n	round/sec	Engine Speed $(\frac{\omega}{2\pi})$

How does dynamical response change due to modeling errors or wear-and-tear?

For example, start time change under probabilistic uncertainty?

**Probabilistic hyperproperties**: Sensitivity under probabilistic parameter change

$$\mathbf{Pr}_{\pi_1,\pi_2}(|\tau^{\pi_1}-\tau^{\pi_2}|\leq \delta)>1-\varepsilon$$

We need new logic to reason over *multiple* random paths!

[Jin et.al, HSCC 14]

## HyperPSTL: Hyper Probabilistic Signal Temporal Logic



STL

Add reference to different paths

HyperSTL

Add probabilistic quantifications

HyperPSTL

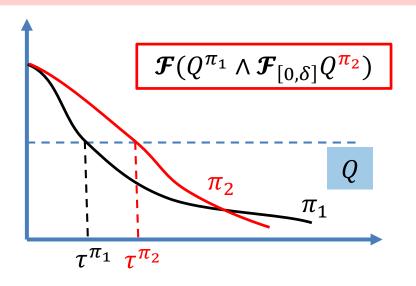
Add probabilistic arithmetic

(full) HyperPSTL

HyperPSTL: 
$$\varphi := \mathbf{a}^{\pi} \mid \varphi^{\pi} \mid \neg \varphi \mid \varphi \wedge \varphi \mid \varphi \mathcal{U}_{[t_1,t_2]} \varphi \mid p \bowtie p$$

$$p := \mathbb{P}^{\Pi} \varphi \mid \mathbb{P}^{\Pi} p \mid c$$

- a ∈ AP, and AP is the finite set of atomic propositions,
- $t_1 < t_2$  with  $t_1, t_2 \in \mathbb{Q}_{\infty}$ ,
- $\pi$  is a path variable, and  $\Pi$  is a set of path variables,
- $\mathbb{P}$  is the probability operator,  $c \in [0,1]$
- ⋈ ∈ {<,>,=,≤,≥},
- $fv(\varphi) = \emptyset$



Probabilistic quantifications of multiple parallel paths

$$\mathbb{P}^{(\pi_1,\pi_2)}\left(\mathcal{F}(Q^{\pi_1}\wedge\mathcal{F}_{[0,\delta]}Q^{\pi_2})\right) > p$$

Nested probabilistic path quantification

$$\mathbb{P}^{\pi_1}\left(\mathbb{P}^{\pi_2}\left(\mathcal{F}\left(Q^{\pi_1}\wedge\mathcal{F}_{[0,\delta]}Q^{\pi_2}\right)\right)>p_2\right)>p_1$$

## HyperPSTL: Hyper Probabilistic Signal Temporal Logic



STL

Add reference to different paths

HyperSTL

Add probabilistic quantifications

HyperPSTL

Add probabilistic arithmetic

(full) HyperPSTL

HyperPSTL: 
$$\varphi \coloneqq \mathbf{a}^{\pi} \mid \varphi^{\pi} \mid \neg \varphi \mid \varphi \land \varphi \mid \varphi \mathcal{U}_{[t_1,t_2]} \varphi \mid p \bowtie p$$
 
$$p \coloneqq \mathbb{P}^{\Pi} \varphi \mid \mathbb{P}^{\Pi} p \mid f(p,...,p)$$

- a ∈ AP, and AP is the finite set of atomic propositions,
- $t_1 < t_2$  with  $t_1, t_2 \in \mathbb{Q}_{\infty}$ ,
- $\pi$  is a path variable, and  $\Pi$  is a set of path
- P is the probability operator,
- ⋈ ∈ {<,>,=,≤,≥},

Kullback-Leibler divergence of two satisfaction probabilities  $\mathbb{P}^{\pi_1} \varphi_1^{\pi_1}$  and  $\mathbb{P}^{\pi_2} \varphi_2^{\pi_2}$ :

$$\mathbb{P}^{\pi_1} \varphi_1^{\pi_1} \log \left( \frac{\mathbb{P}^{\pi_1} \varphi_1^{\pi_1}}{\mathbb{P}^{\pi_2} \varphi_2^{\pi_2}} \right) + \left( 1 - \mathbb{P}^{\pi_1} \varphi_1^{\pi_1} \right) \log \left( \frac{1 - \mathbb{P}^{\pi_1} \varphi_1^{\pi_1}}{1 - \mathbb{P}^{\pi_2} \varphi_2^{\pi_2}} \right) < c$$

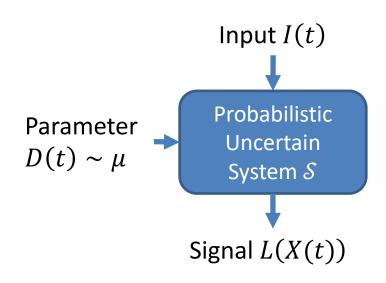
- $f: \mathbb{R}^n \to \mathbb{R}$  is a n-ary elementary function, constants are viewed as 0-ary functions,
- $fv(\varphi) = \emptyset$

## Semantics: HyperPSTL on Probabilistic Uncertain System



#### **Probabilistic uncertain system** (PUS): $S = (X, I, D, \mu, AP, L)$ where

- $\mathcal{X}$  is the state space,  $X^{\text{init}} = (x_1^{\text{init}}, ..., x_l^{\text{init}}) \in \mathcal{X}$  is an initial state
- Parameter  $D(t)=(d_1(t),\ldots,d_n(t))$  for  $t\in\mathbb{R}_{\geq 0}$  is drawn from probability distribution  $\mu$
- Input  $I(t) = (i_1(t), ..., i_m(t))$  is an m-dimensional function of time t
- Given I(t) and D(t), the system generates a **path**  $X: \mathbb{R}_{\geq 0} \to \mathcal{X}$  with  $X(t) = (x_1(t), ..., x_l(t))$
- AP is a set of atomic propositions, L:  $\mathcal{X} \to 2^{\mathsf{AP}}$  is a labeling function
- a path of the system induces a signal  $\sigma(t) = L(X(t))$ :  $\mathbb{R}_{\geq 0} \to 2^{AP}$ .



#### The PUS modeling allows capturing

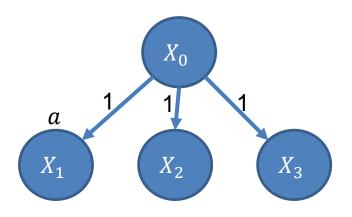
- Hybrid I/O automata with probabilistic parameters (e.g., powertrain)
- continuous-time Markov chains (CTMCs) as in queueing networks

## **HyperPSTL: Expressiveness**



**Theorem**: HyperPSTL *strictly subsumes* PSTL (its non-hyper fraction) on CTMCs.

Prof idea: find a CTMC and a property, such that this property only be expressed in HyperPSTL



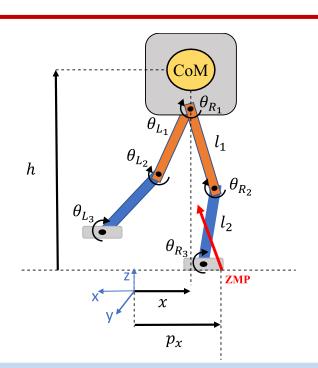
- CTMC has only 3 paths
- Satisfaction probability of any STL is  $0, \frac{1}{3}, \frac{2}{3}, 1$ , so  $P(\varphi) = \frac{1}{9}$  is always false for any  $(\varphi)$
- HyperPSTL  $P^{(\pi_1,\pi_2)}(\mathcal{F}(a^{\pi_1} \wedge a^{\pi_2})) = \frac{1}{9}$  is true

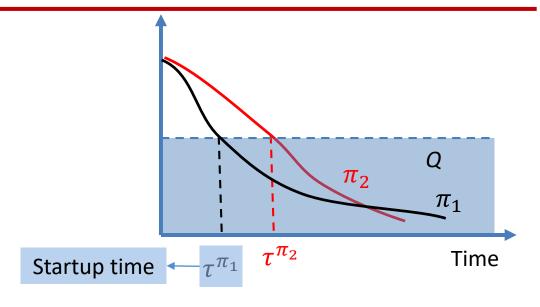
## HyperPSTL in Action: Sensitivity to Modeling Errors



#### **Toyota Powertrain Benchmark**







Dynamical response depends on system parameters

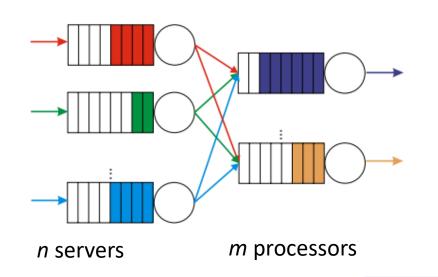
**Design specification**: Sensitivity of startup time

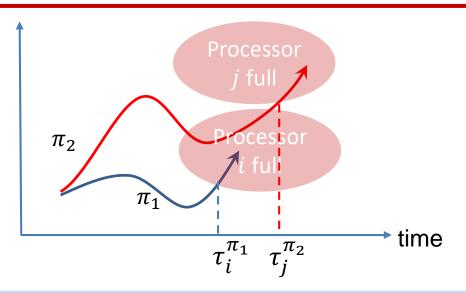
$$\mathbf{Pr}_{\pi_1,\pi_2}(|\tau^{\pi_1}-\tau^{\pi_2}|\leq \delta)>1-\varepsilon$$

$$\mathbb{P}^{(\pi_1,\pi_2)} \begin{pmatrix} (\neg Q^{\pi_1} \wedge \neg Q^{\pi_2}) \\ \boldsymbol{u} \left( (Q^{\pi_1} \wedge \boldsymbol{\mathcal{F}}_{[0,\delta]} Q^{\pi_2}) \vee (Q^{\pi_2} \wedge \boldsymbol{\mathcal{F}}_{[0,\delta]} Q^{\pi_1}) \right) \end{pmatrix} > 1 - \varepsilon$$

## **HyperPSTL** in Action: Workload Fairness







#### **Design specification**: Workload Fairness

$$\mathbf{Pr}_{\pi_{1}}\left(\left|\mathbf{Pr}_{\pi_{2}}\left(\tau_{i}^{\pi_{1}}-\tau_{j}^{\pi_{2}}>t\right)-\mathbf{Pr}_{\pi_{2}}\left(\tau_{i}^{\pi_{1}}-\tau_{j}^{\pi_{2}}>t\right)\right|<\delta\right)>1-\varepsilon$$



This should hold with probability more than  $1-\varepsilon$  for  $\pi_1$ 



For any fixed  $\tau_i^{\pi_1}$ , the probability difference between  $\tau_i^{\pi_1} - \tau_j^{\pi_2} > \mathbf{t}$  and  $\tau_i^{\pi_1} - \tau_j^{\pi_2} < -\mathbf{t}$  should be less than  $\delta$ 

$$\mathbb{P}^{\pi_1}(|\mathbb{P}^{\pi_2}((\neg Q_i^{\pi_1} \land \neg Q_j^{\pi_2})\mathcal{U}(Q_i^{\pi_1} \land \diamondsuit_{[\tau,\infty)}Q_j^{\pi_2})) - \mathbb{P}^{\pi_2}((\neg Q_i^{\pi_1} \land \neg Q_j^{\pi_2})\mathcal{U}(Q_j^{\pi_2} \land \diamondsuit_{[\tau,\infty)}Q_i^{\pi_1}))| \leq \delta) \geq 1 - \varepsilon.$$

## HyperPSTL in Action: Probabilistic Detectability



Captured independently of the type of used **sound** detector as probabilistic **overshoot observability** on system outputs, when input overshoot captures that an anomaly has occurred

• Let x be the input and y be the output. After a "step" event, the output signal should be different if the input (1) stays bounded or (2) overshoots.

$$\mathbb{P}^{\{\pi,\pi'\}}((\Box(step^{\pi}\Rightarrow \Box_{I}(x^{\pi}< c)) \land \diamondsuit(step^{\pi'} \land \diamondsuit_{I}(x^{\pi'}> c))) \Rightarrow (\diamondsuit_{I} d(y^{\pi}, y^{\pi'}) > c')) > 1 - \varepsilon$$

## **SMC** of HyperPSTL: Overview



Hyper features beyond existing methods for Statistical Model Checking (SMC)

- Probabilistic quantifications of *multiple* parallel paths (e.g., sensitivity)  $\mathbb{P}^{(\pi_1,\pi_2)} \varphi^{(\pi_1,\pi_2)} < p$
- Nested probabilistic path quantification (e.g., fairness)  $\mathbb{P}^{\pi_1}(\mathbb{P}^{\pi_2}\varphi^{(\pi_1,\pi_2)} < p_2) < p_1$
- **Joint** probabilities (e,g., KL-divergence)  $(\mathbb{P}^{\Pi_1}\varphi_1, \mathbb{P}^{\Pi_2}\varphi_2) \in D$

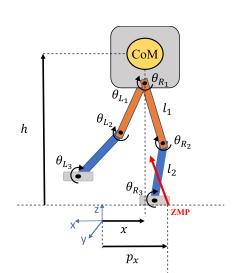
$$\mathbb{P}^{\pi_1} \varphi_1^{\pi_1} \log \left( \frac{\mathbb{P}^{\pi_1} \varphi_1^{\pi_1}}{\mathbb{P}^{\pi_2} \varphi_2^{\pi_2}} \right) + \left( 1 - \mathbb{P}^{\pi_1} \varphi_1^{\pi_1} \right) \log \left( \frac{1 - \mathbb{P}^{\pi_1} \varphi_1^{\pi_1}}{1 - \mathbb{P}^{\pi_2} \varphi_2^{\pi_2}} \right) < c$$



$$\left(\mathbb{P}^{\pi_1} \varphi_1^{\pi_1}, \mathbb{P}^{\pi_2} \varphi_2^{\pi_2}\right) \in D \text{ with } D = \left\{ (x_1, x_2) \mid x_1 \log \left(\frac{x_1}{x_2}\right) + (1 - x_1) \log \left(\frac{1 - x_1}{1 - x_2}\right) < c \right\}$$

### **Sensitivity Verification of Real-World CPS**





$$\mathbb{P}^{(\pi_1,\pi_2)} \begin{pmatrix} (\neg Q^{\pi_1} \wedge \neg Q^{\pi_2}) \\ u \left( (Q^{\pi_1} \wedge \mathcal{F}_{[0,\delta]} Q^{\pi_2}) \vee (Q^{\pi_2} \wedge \mathcal{F}_{[0,\delta]} Q^{\pi_1}) \right) \end{pmatrix} > 1 - \varepsilon$$

## Walking Robot Benchmark With Reinforcement Learning Controller

δ	ε	α	Acc.	Sam.	Time (s)	Ans.
2.4	0.02	0.01	1.00	7.4e+01	3.0e-01	False
2.4	0.02	0.05	0.99	4.4e+01	1.4e-01	False
2.4	0.12	0.01	1.00	4.2e+01	1.2e-01	True
2.4	0.12	0.05	1.00	2.1e+01	7.0e-02	True
2.4	0.2	0.01	1.00	1.3e+01	4.0e-02	True
3.0	0.02	0.01	1.00	1.1e+01	2.4e-02	False
3.0	0.02	0.05	1.00	6.5e+00	1.1e-02	False
3.0	0.12	0.05	0.98	7.0e+01	4.3e-01	False
3.0	0.2	0.01	1.00	1.6e+02	5.5e-01	True
3.0	0.2	0.05	0.98	1.0e+02	2.9e-01	True

#### **Toyota Powertrain Benchmark**

δ	ε	α	Acc.	Sam.	Time (s)	Ans.
0.1	5 0.95	0.05	1.00	5.9e+01	8.1e+00	True
0.1	5 0.95	0.01	1.00	9.0e+01	1.3e+01	True
0.1	5 0.99	0.05	0.99	6.6e+01	9.1e+00	False
0.1	5 0.99	0.01	1.00	9.7e+01	1.4e+01	False
0.2	0 0.95	0.05	0.98	5.9e+01	8.1e+00	True
0.2	0 0.95	0.01	1.00	9.0e+01	1.2e+01	True
0.2	0 0.99	0.05	1.00	3.0e+02	4.2e+01	True
0.2	0 0.99	0.01	0.99	4.6e+02	1.8e+02	True

### [EMSOFT'19, TACAS'20\*]



 On continuous-time probabilistic models (e.g., powertrain, queueing network), how to capture properties between many paths (sensitivity, fairness, attack detectability)?



Hyper Probabilistic Signal Temporal Logics: *HyperPSTL* 

How to reason about HyperPSTL on complex systems?



Statistical Model Checking (*SMC*) of HyperPSTL

How does the SMC work in practice?



Evaluation on real-world CPS

#### **Current Work:**

- Application to conformance testing
- Synthesis by reinforcement learning

## Thank you











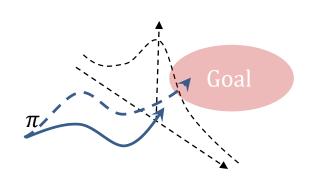




### What is a probabilistic hyperproperty?

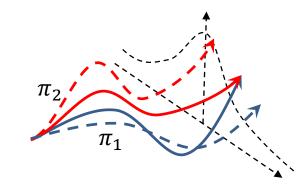


#### Probabilistic hyperproperty reasons over **multiple** random paths.



# Probabilistic Property:

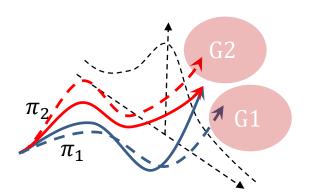
• Reachability  $Pr(\pi \models \mathcal{F}(Goal))$  > 0.99



#### Probabilistic Hyperproperty:

Two path meet

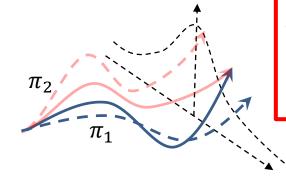
$$\Pr((\pi_1, \pi_2) \models \mathcal{F}(\pi_1 = \pi_2))$$
  
> 0.99



#### Probabilistic Hyperproperty:

 Compare satisfaction probabilities

$$Pr(\pi \models \mathcal{F}(G1))$$
  
>  $Pr(\pi \models \mathcal{F}(G2))$ 



• One catchup another  $\mathbf{Pr}(\pi_1 \models \mathbf{C}) > 0.5$ 

where

$$C: \Pr(\pi_2 \models \mathcal{F}(\pi_1 = \pi_2))$$
  
> 0.99