

Exploitation of Rotational Symmetries to Solve the Swarm Initialization Problem

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- **We have developed a sound, mathematical basis for solving the swarm initialization problem in a special case.**
 - We will assume a circular swarm trajectory (eccentricity = 0, exact)
 - We define two operations which preserve swarm optimality:
 - Rotation of the swarm within some known space of valid rotations;
 - Transposition of any two satellites at any point in the orbit.
- **Quantization of the solution space** has enabled the use of discretized optimization methods – chiefly, the *Munkres Algorithm*.
- **A staged optimization approach** solves the problem in successively more detailed passes.
- **Computational slowness** is the current major obstacle to implementing this method.
 - Careful algorithm selection;
 - Parallelization of the process;
 - Dividing stages by regularity of use;

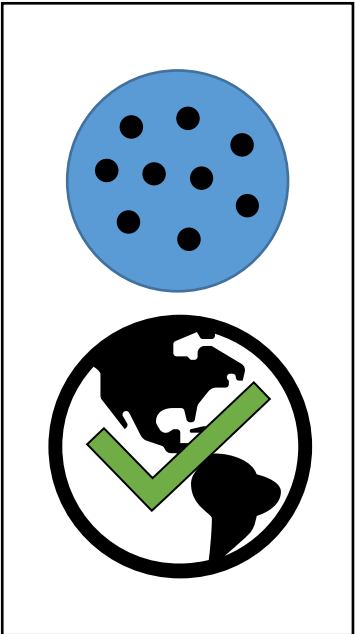


In prior discussions, a swarm was defined to be a close-flying formation of satellites for which the following quantities could be defined:

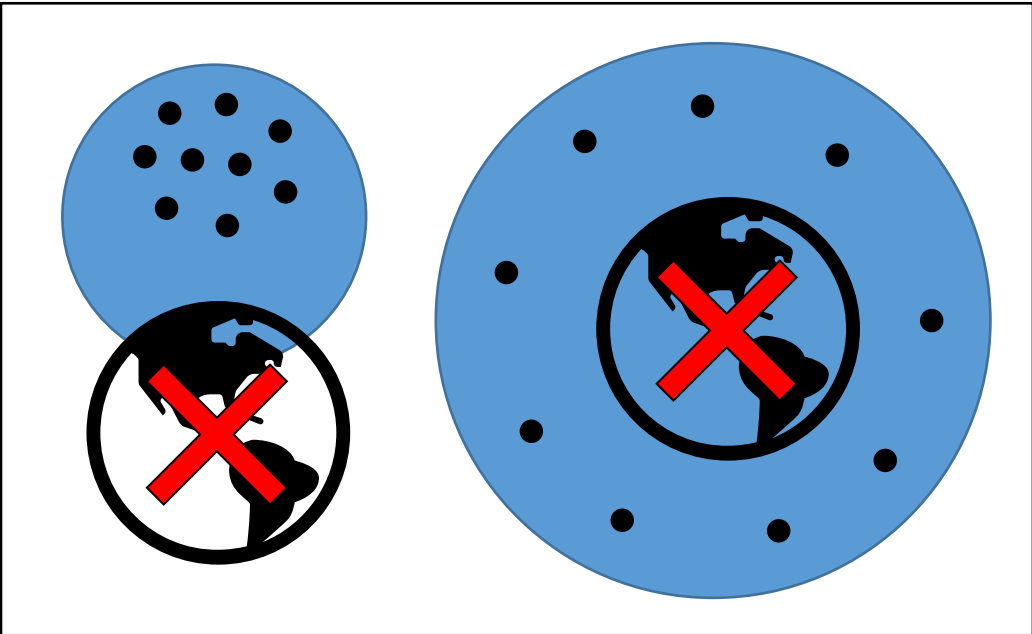
- **Swarm envelope**
 - A **closed, convex set of points with no holes, gaps, or voids**, which contains all satellites in the swarm.
 - Mathematically speaking, **the boundary of the envelope must be simply connected.**
 - The swarm envelope **may not contain points inaccessible to the satellites** (e.g., points at or below Earth's surface).
 - **For this discussion, we will use a spherical envelope.**
 - Must include a **reference point** that marks the envelope center.
 - This point **need not to be an element** of the swarm envelope.
 - **For this discussion, we will use the center of the sphere.**



Principal Definitions



VALID



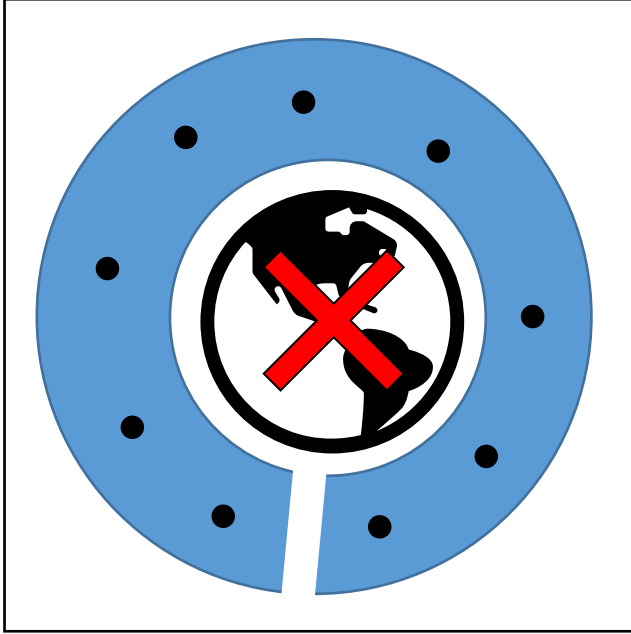
CONTAIN INACCESSIBLE POINTS



Principal Definitions



CONTAINS A VOID



NOT CONVEX

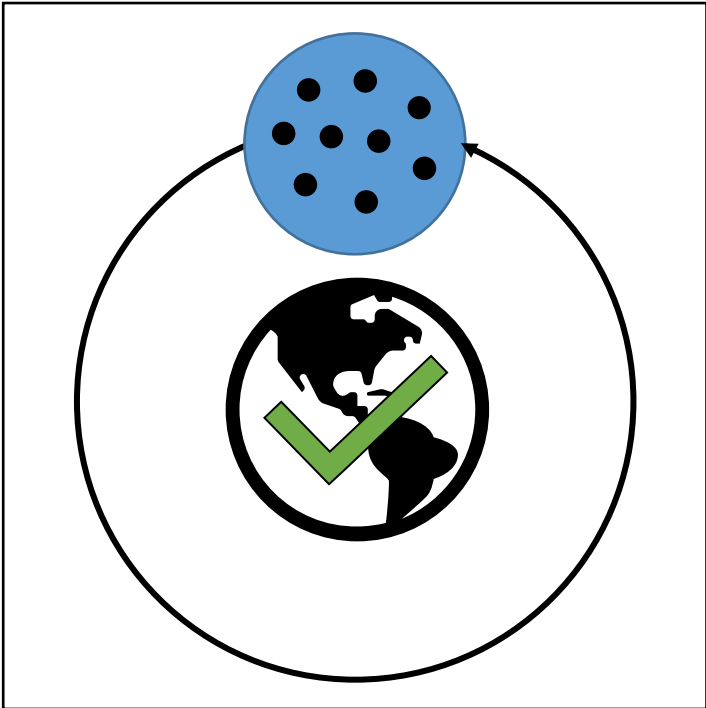


- **Swarm trajectory**

- A **user-defined function of time** which specifies the position of the swarm envelope relative to any point with known coordinates.
 - This latter point can be taken to be the **center of the Earth**.
 - Position measured to the **envelope-relative reference point**.
- If necessary, this definition may include a function to specify **the orientation of the swarm envelope**.
 - Unnecessary in the case of a spherical envelope.
- While it is not a strict requirement in the general case, **trajectories should follow spacetime geodesics**.
- For this discussion, we will assume that the swarm trajectory **can be approximated by a Keplerian orbit of zero eccentricity**.



Principal Definitions



FOLLOWS A
GEODESIC PATH



FOLLOWS ONLY
ITS DREAMS



- **Swarm distribution**

- The **arrangement of satellites** within the sphere envelope as defined relative to the envelope-relative reference point.
- The most abstract component of the satellite swarm, incorporating mission-specific parameters, an unknown number of degrees of freedom, and the swarm cost functional.
- For this discussion, we will assume that the swarm distribution **can be rotated** about the envelope relative reference point.

- **Swarm centroid**

- A function of the satellite positions which **determines the true center of the swarm**.
- Distinct from the envelope-relative reference point, as the centroid is **determined by satellite positions, not envelope geometry**.
- For this discussion, we will use the **arithmetic mean** of all satellite positions to define the swarm centroid.



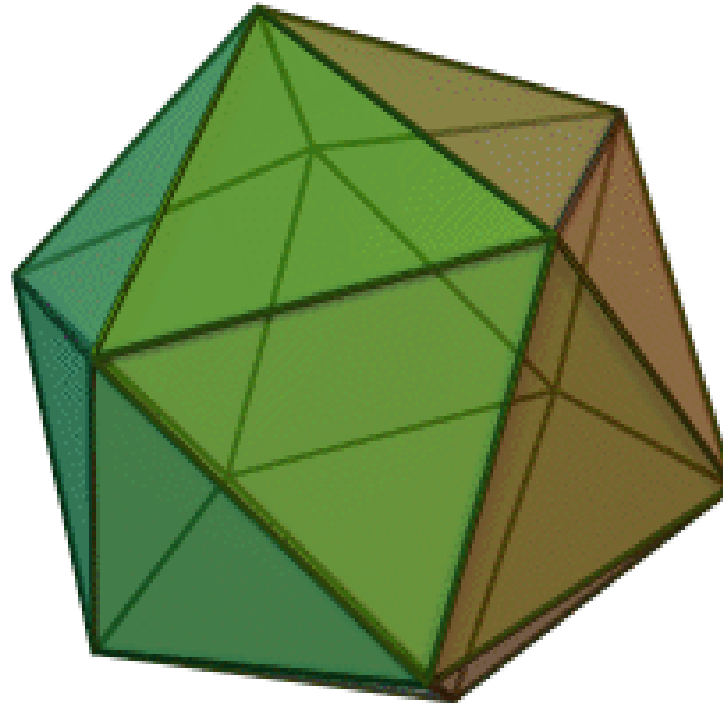
For this discussion, we will default to our test case of **twelve satellites**:

- The swarm envelope is a **sphere of radius ρ** with reference at its center.
- The swarm trajectory is a **circular orbit of radius a** .
- The satellites in the swarm should be distributed at the **vertices of a regular icosahedron**.
 - The swarm may be **transformed by any rotation** about its center.
- The swarm centroid is defined using the arithmetic mean:

$$\vec{r}_0(t) = \frac{1}{12} \sum_{i=1}^{12} \vec{r}_i(t)$$



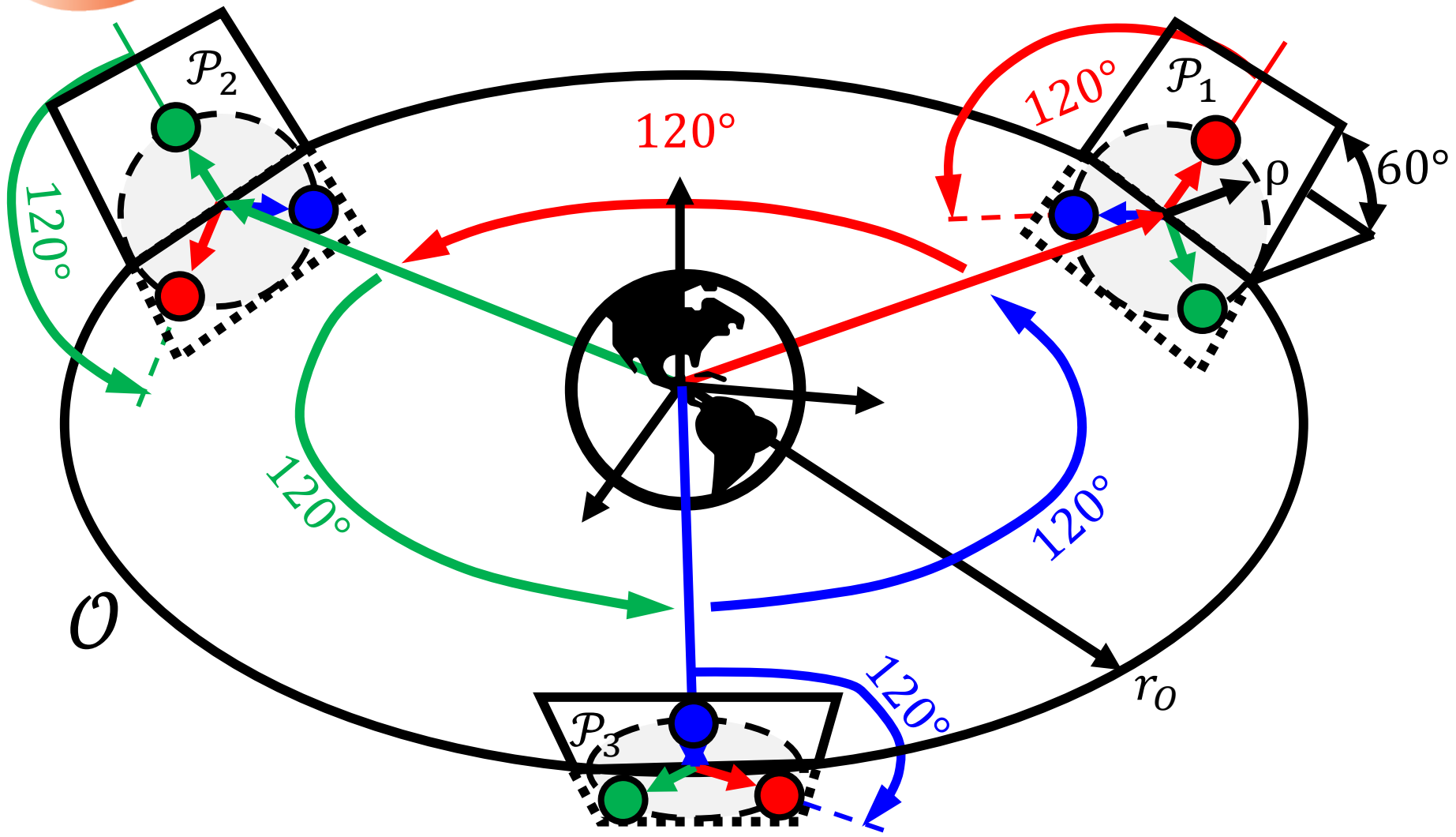
Test Swarm Geometry



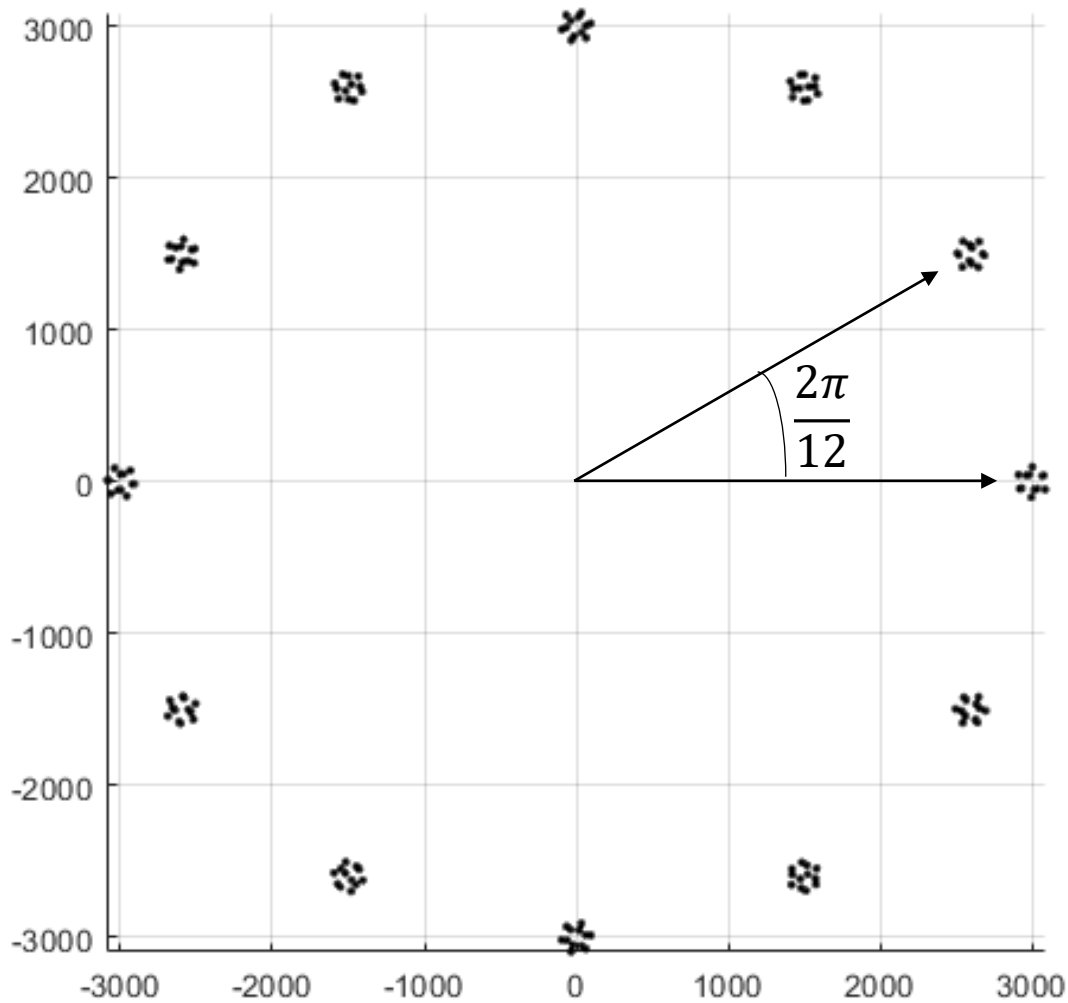
A rotating icosahedron is still an icosahedron.

Animation from Wikipedia.

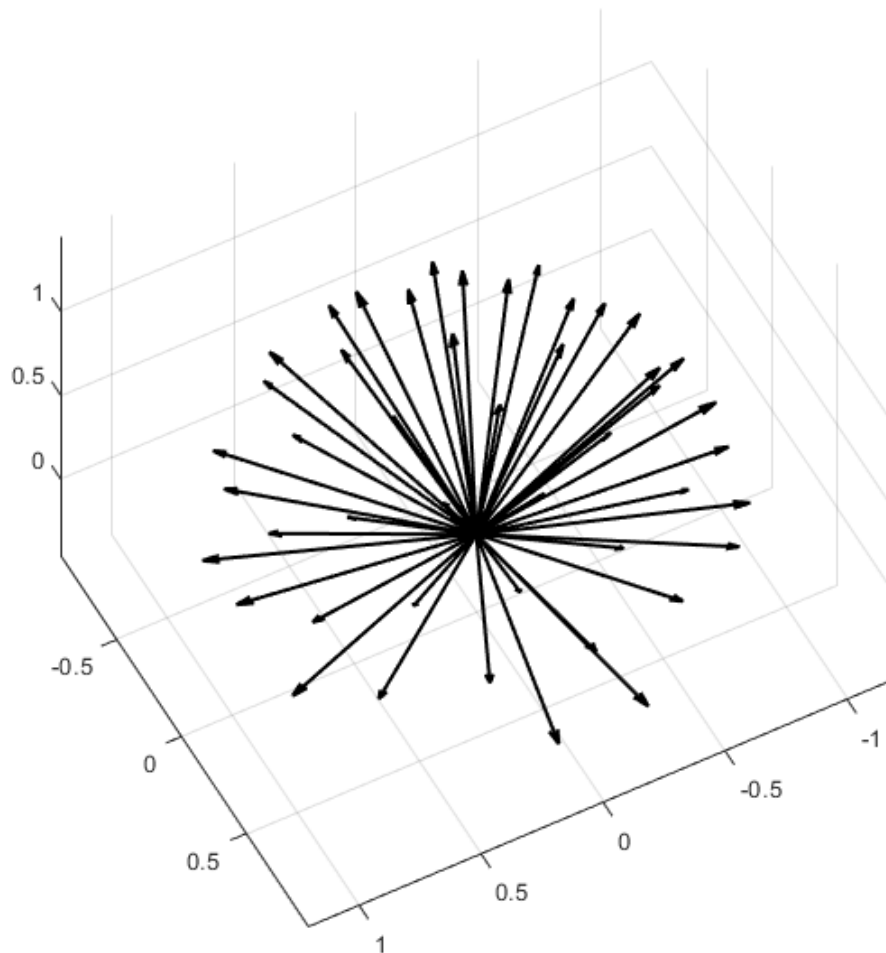
Test Swarm Configuration



Formation Chains



Formation Chains





- Note that n successive rotations by $2\pi/n$ radians about a given axis will result in a net angle of 2π radians.
- We may, however, rotate by *multiples* of $2\pi/n$ radians.
 - Recognizing that a rotation by 2π radians is equivalent to no rotation, it follows that rotations by multiples of $n + 1$ to $2n - 1$ are equivalent to rotations by **multiples of 1 to $n - 1$** .
 - We may also consider the negative multiples (i.e., $-(n - 1)$ to -1).
 - Rotations by a multiple of 0 need only be considered about a single axis, since this corresponds to zero rotation.
- With two distinct rotations to check (the starting attitude and the change in attitude), we must check **$4m^2n(n - 1) + 1$ formation chains**.
- **For 46 axes and 12 satellites, this corresponds to 1,117,249 chains.**
- Some of these chains may be sufficiently close to others to qualify as “duplicate.” **We may eliminate these duplicate chains.**
 - This process takes several days to a week or more on a single CPU, but only needs to be done once (the results may be saved and re-used).

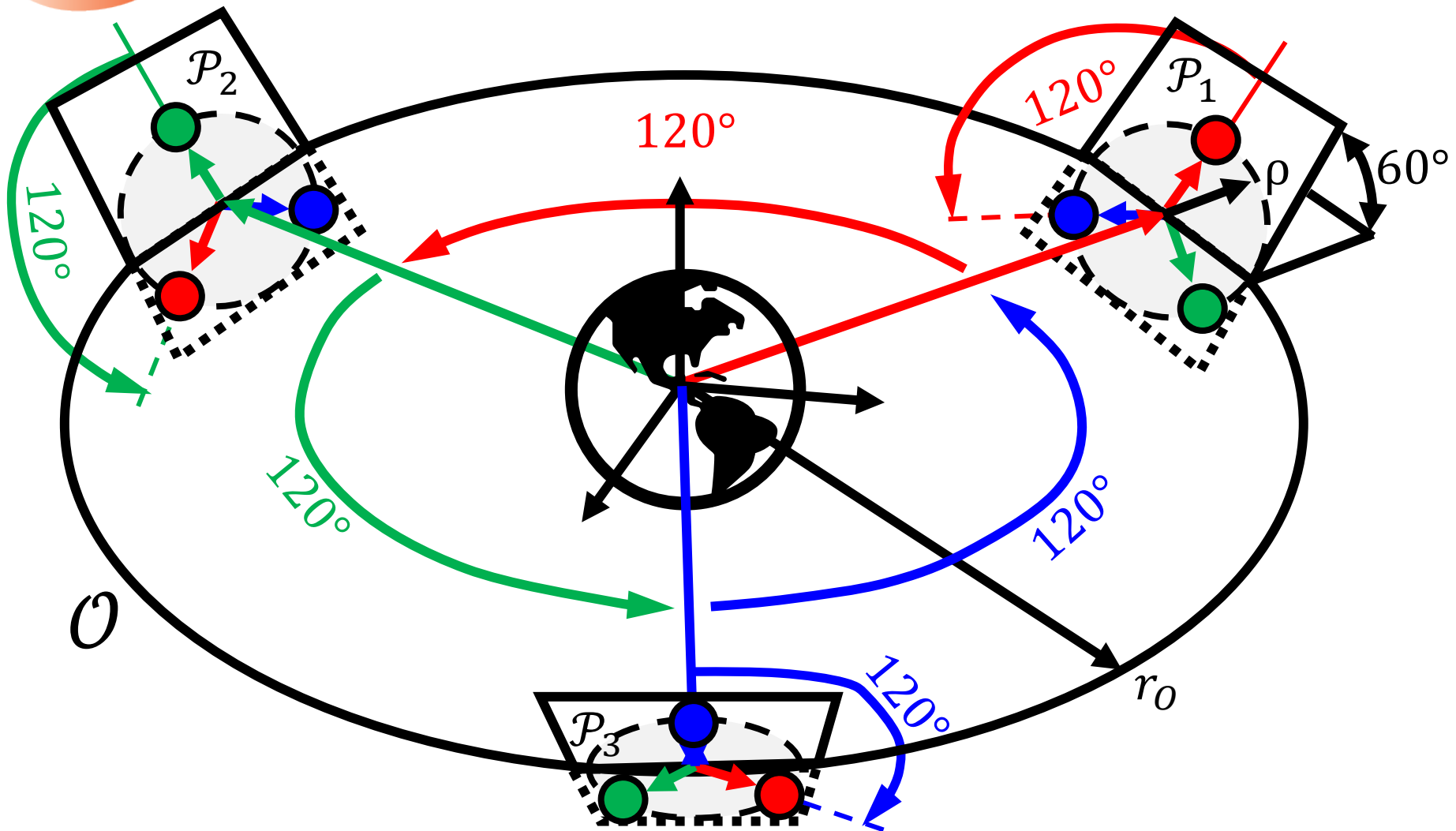


- For 46 axes and 12 satellites, **with duplicate elimination**, we obtain **between 5,000 and 10,000** formation chains depending on tolerance.
- **These formation chains are saved** and used in the next step.
 - Note that, because these chains depend only on the selected value of the rotation, this process is **invariant with respect to the radii of the formation, or its orbit.**



- Once a group of formation chains are selected, we must then pull orbits from them by **selecting the first two formations out of the chain.**
- For this pair, we may connect **each satellite from the first formation to any position from the second.**
 - Connecting the i th satellite to the j th position yields an orbit. Comparing this orbit to the formation chain yields a contribution to a cost function, J_{ij} .
 - For each i and j , we select the two formations with minimal J_{ij} .
- Note that, for each pair, we may propagate the assignment to another formation, so we must compare the values of J_{ij} for each formation in the chain.
- The optimal assignment may be determined using the Munkres Algorithm

Selecting Orbits





- The Munkres algorithm runs in $O(n^4)$ time, so we are currently investigating ways to avoid running the Munkres algorithm in cases where it would clearly produce a suboptimal result.
 - For example, if the sum of the n smallest elements in the cost matrix, J is greater than the current best result, we may skip the algorithm.
- Once the best result has been found, that formation chain is selected, along with the assignment function determined by Munkres and the orbits it generates.
- Those orbits are passed on to the next stage of optimization, which uses continuous optimization to further refine the initial velocities and minimize the swarm cost functional.



- Implement code to run the orbit selection protocol.
- Implement code to refining stage.
- Run this with known test cases, such as the LISA configuration.
- Compare performance for low- n vs. high- n cases.