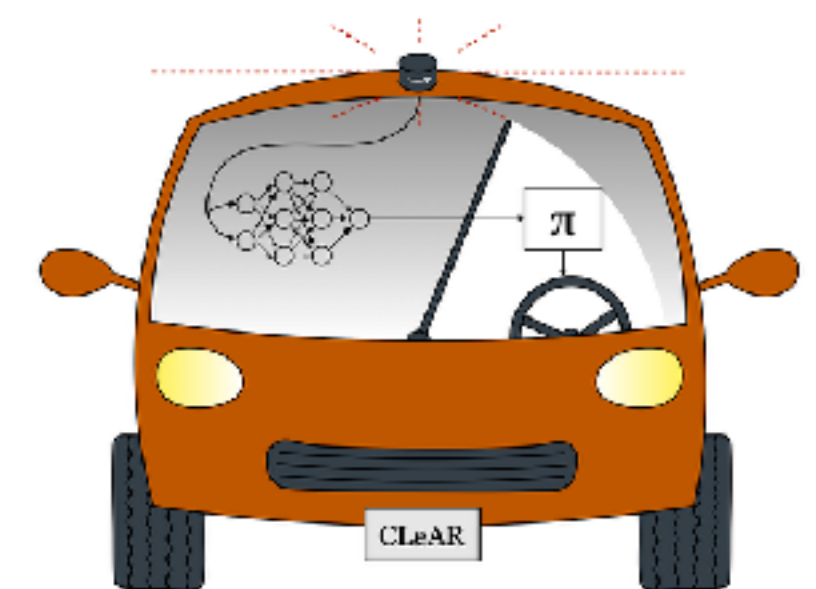


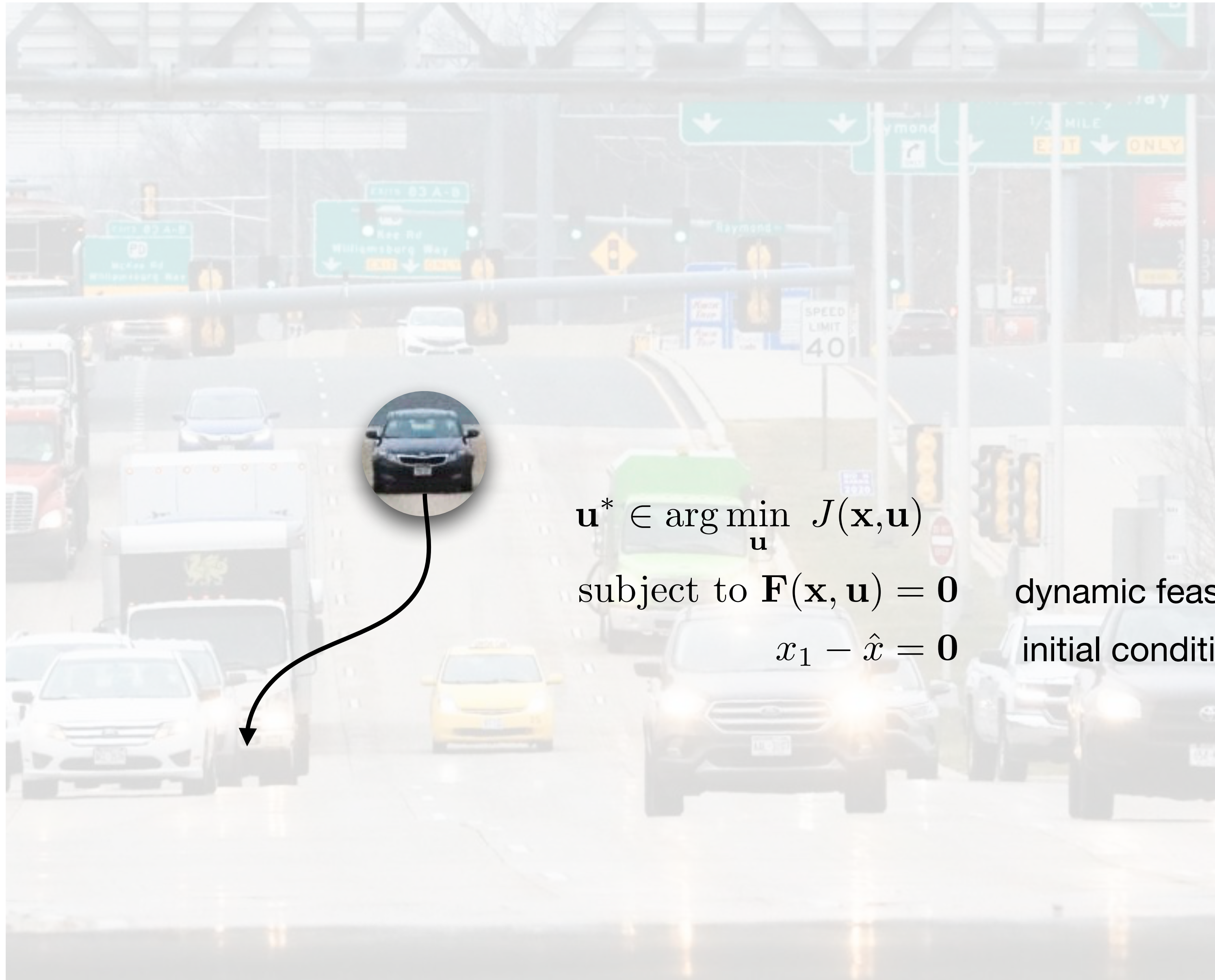
Mixing Continuous Strategies: A Case Study in Trajectory Games

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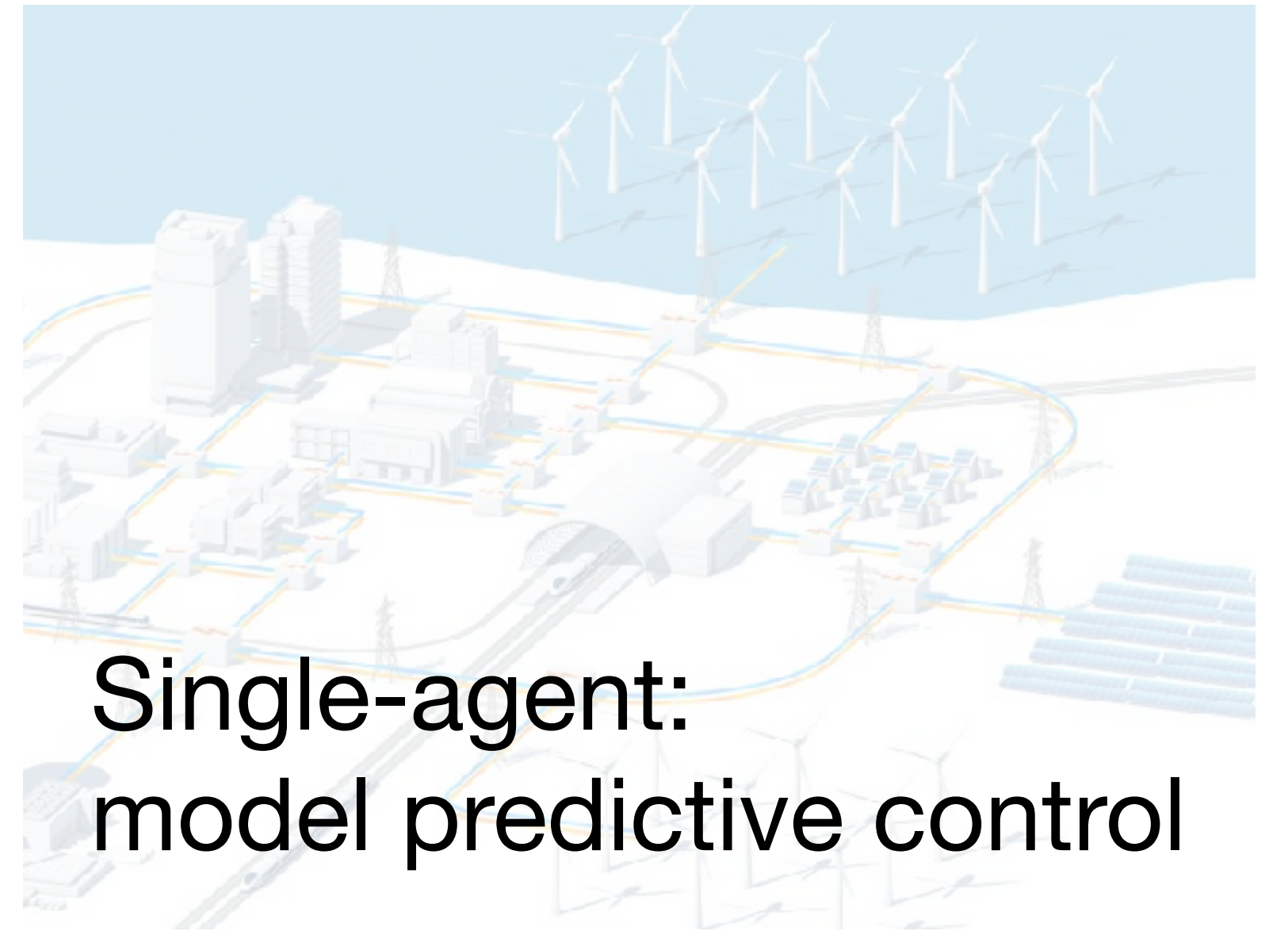




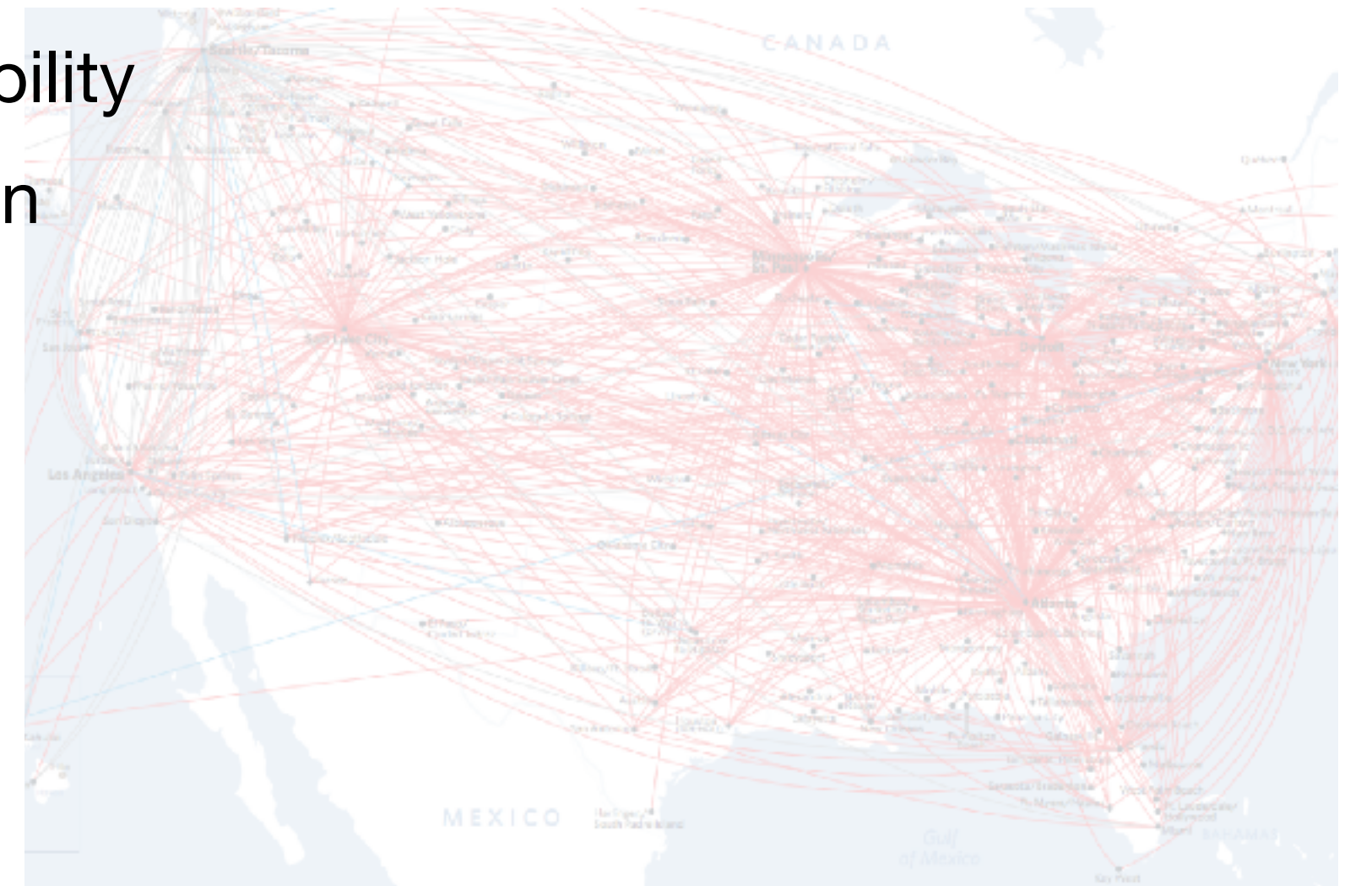
$$\mathbf{u}^* \in \arg \min_{\mathbf{u}} J(\mathbf{x}, \mathbf{u})$$

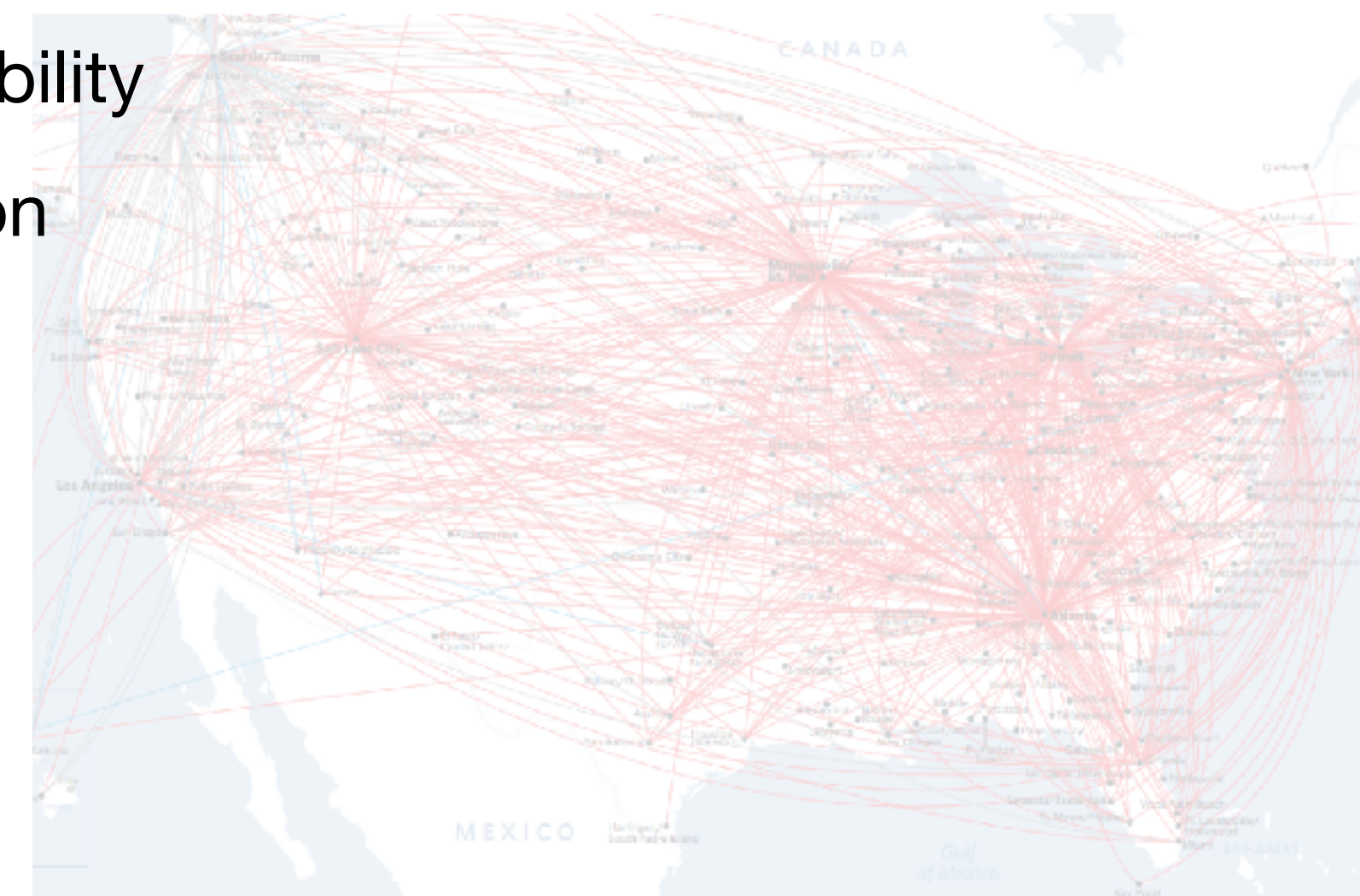
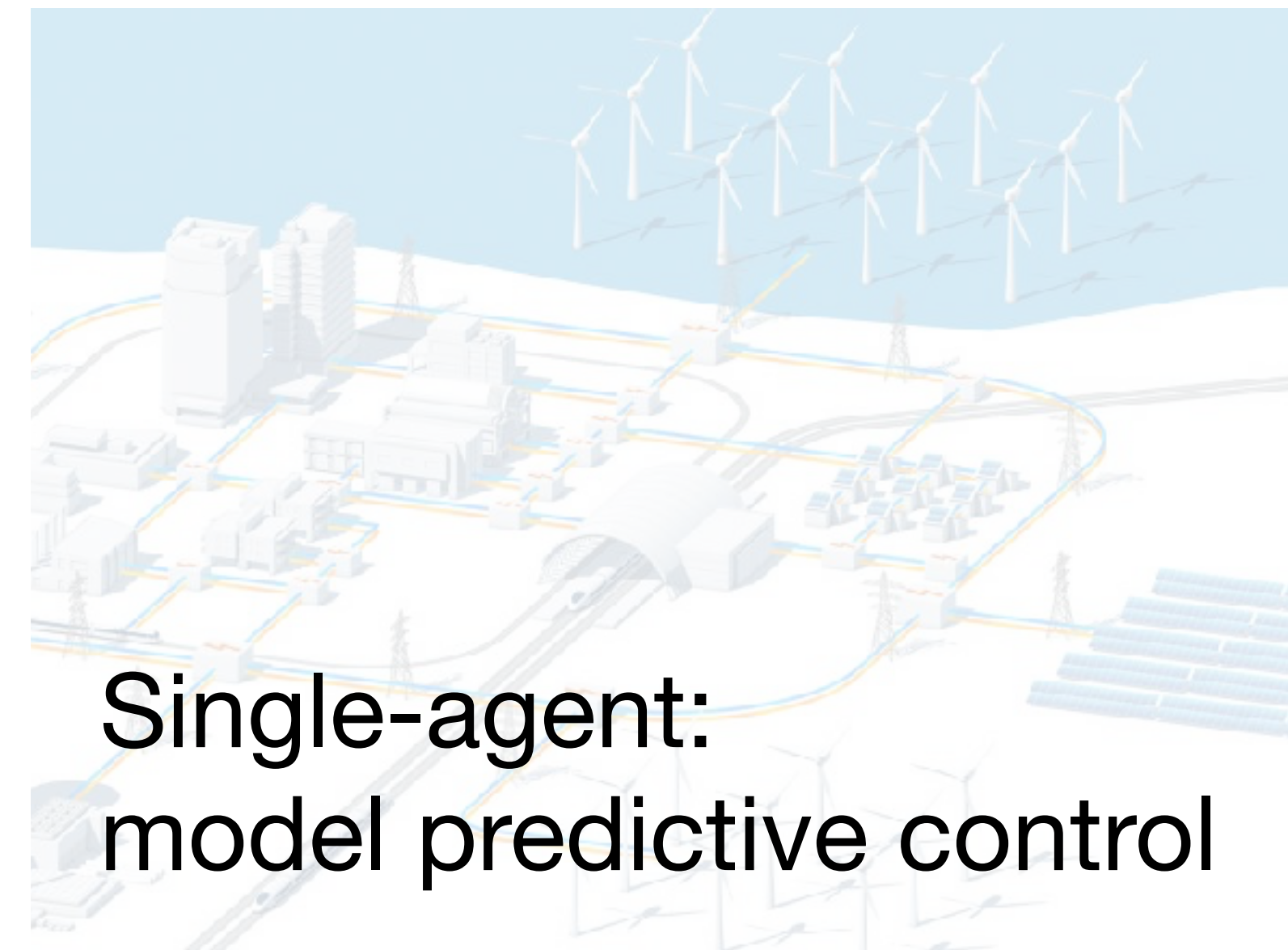
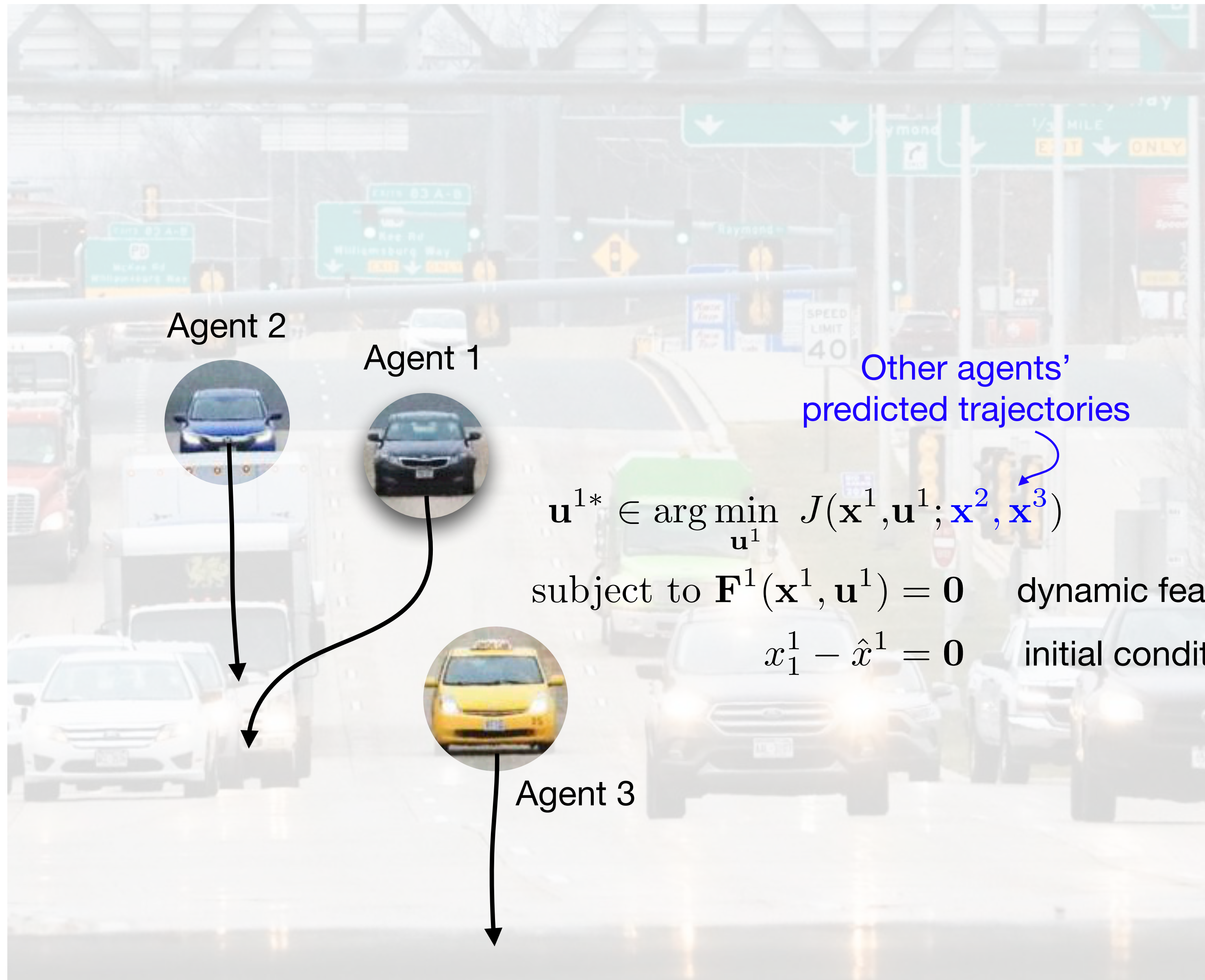
subject to $\mathbf{F}(\mathbf{x}, \mathbf{u}) = \mathbf{0}$ dynamic feasibility

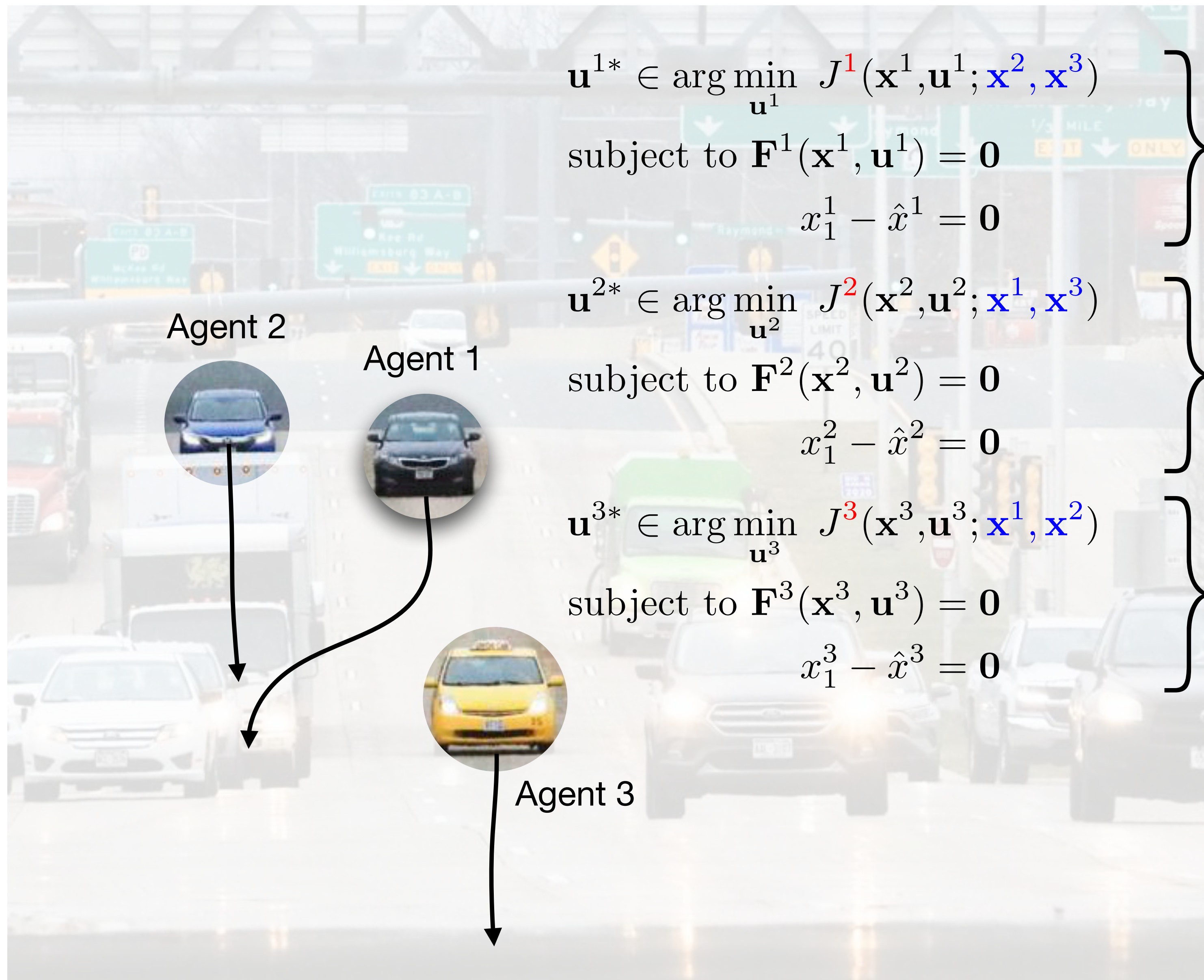
$x_1 - \hat{x} = 0$ initial condition



Single-agent:
model predictive control







$$\mathbf{u}^{1*} \in \arg \min_{\mathbf{u}^1} J^1(\mathbf{x}^1, \mathbf{u}^1; \mathbf{x}^2, \mathbf{x}^3)$$

$$\text{subject to } \mathbf{F}^1(\mathbf{x}^1, \mathbf{u}^1) = \mathbf{0}$$

$$x_1^1 - \hat{x}^1 = 0$$

$$\mathbf{u}^{2*} \in \arg \min_{\mathbf{u}^2} J^2(\mathbf{x}^2, \mathbf{u}^2; \mathbf{x}^1, \mathbf{x}^3)$$

$$\text{subject to } \mathbf{F}^2(\mathbf{x}^2, \mathbf{u}^2) = \mathbf{0}$$

$$x_1^2 - \hat{x}^2 = 0$$

$$\mathbf{u}^{3*} \in \arg \min_{\mathbf{u}^3} J^3(\mathbf{x}^3, \mathbf{u}^3; \mathbf{x}^1, \mathbf{x}^2)$$

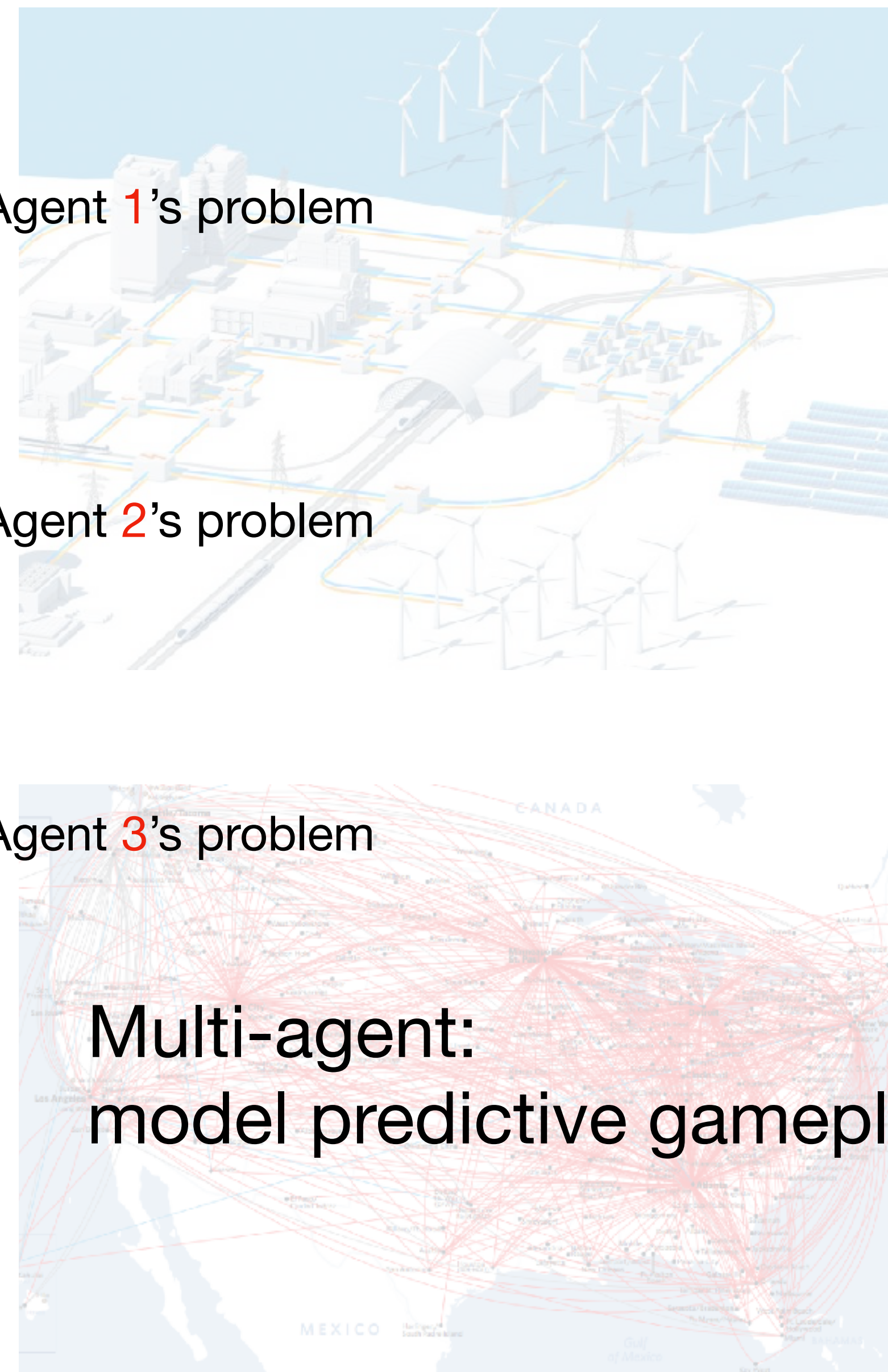
$$\text{subject to } \mathbf{F}^3(\mathbf{x}^3, \mathbf{u}^3) = \mathbf{0}$$

$$x_1^3 - \hat{x}^3 = 0$$

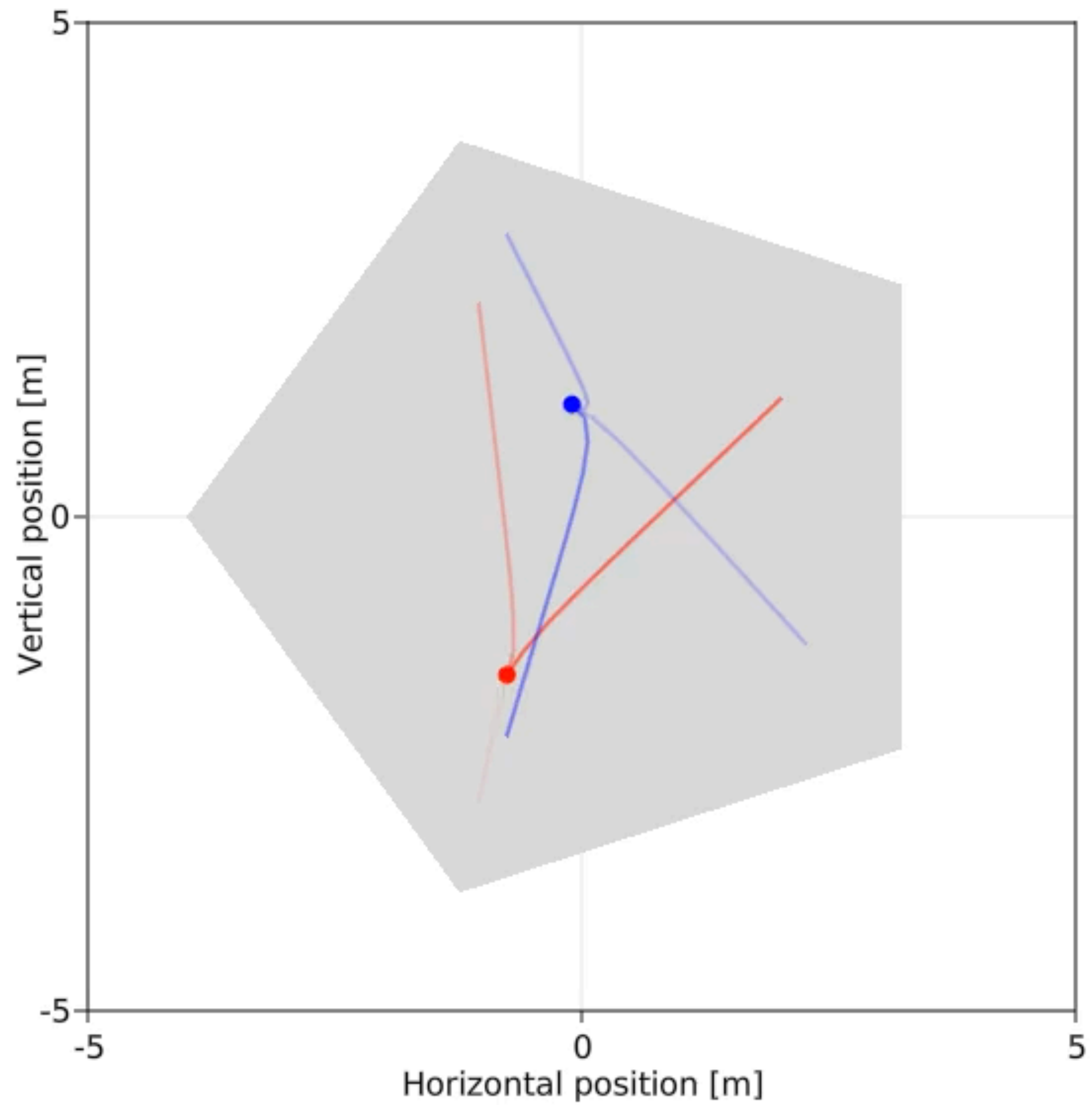
Agent 1's problem

Agent 2's problem

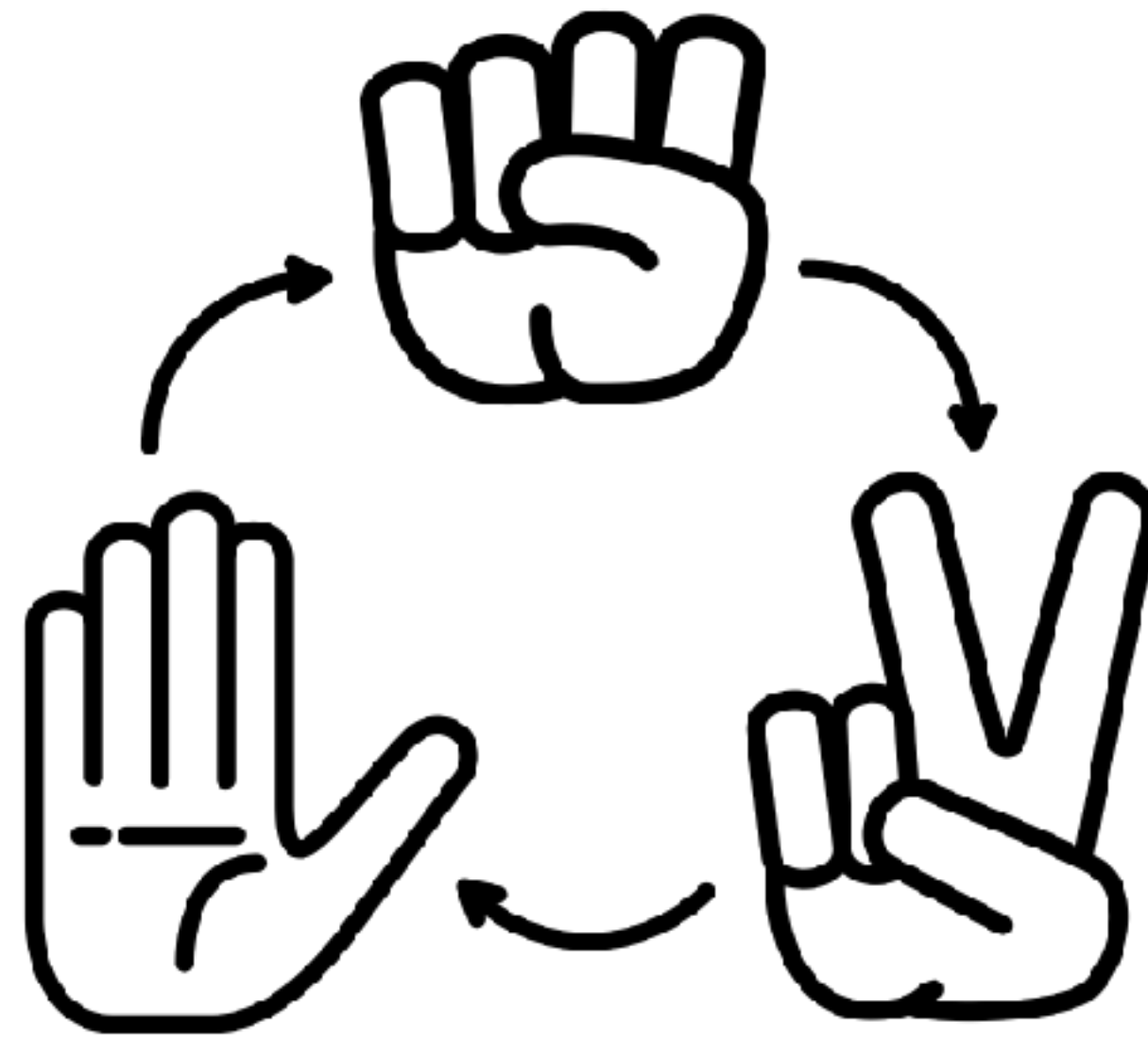
Agent 3's problem



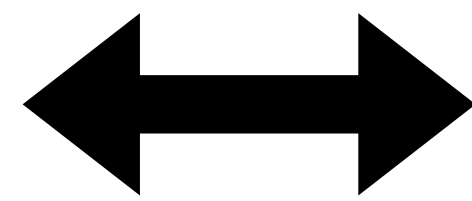
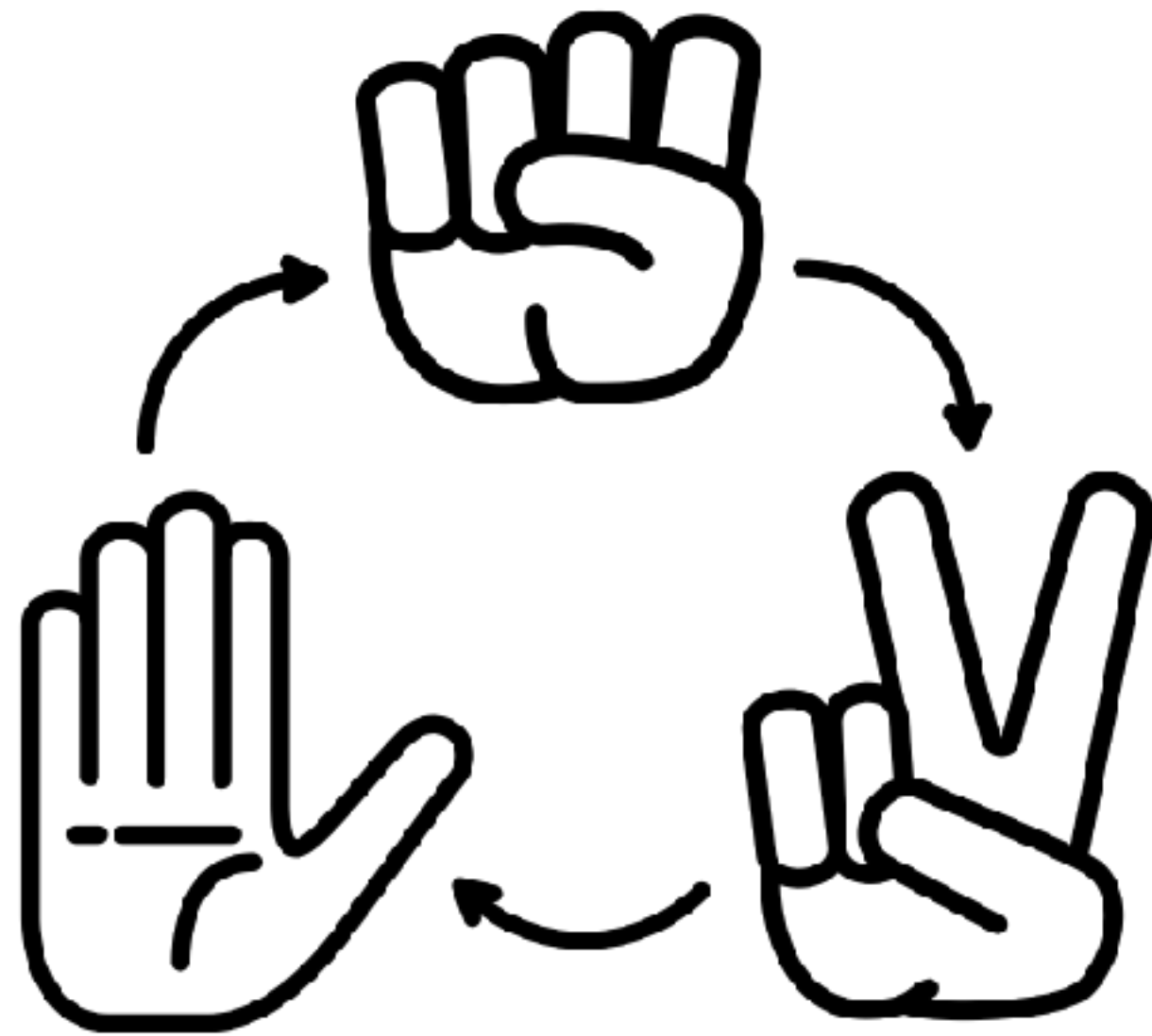
Multi-agent:
model predictive gameplay



“Pure” Nash equilibria may not exist

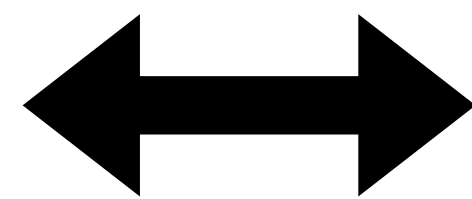
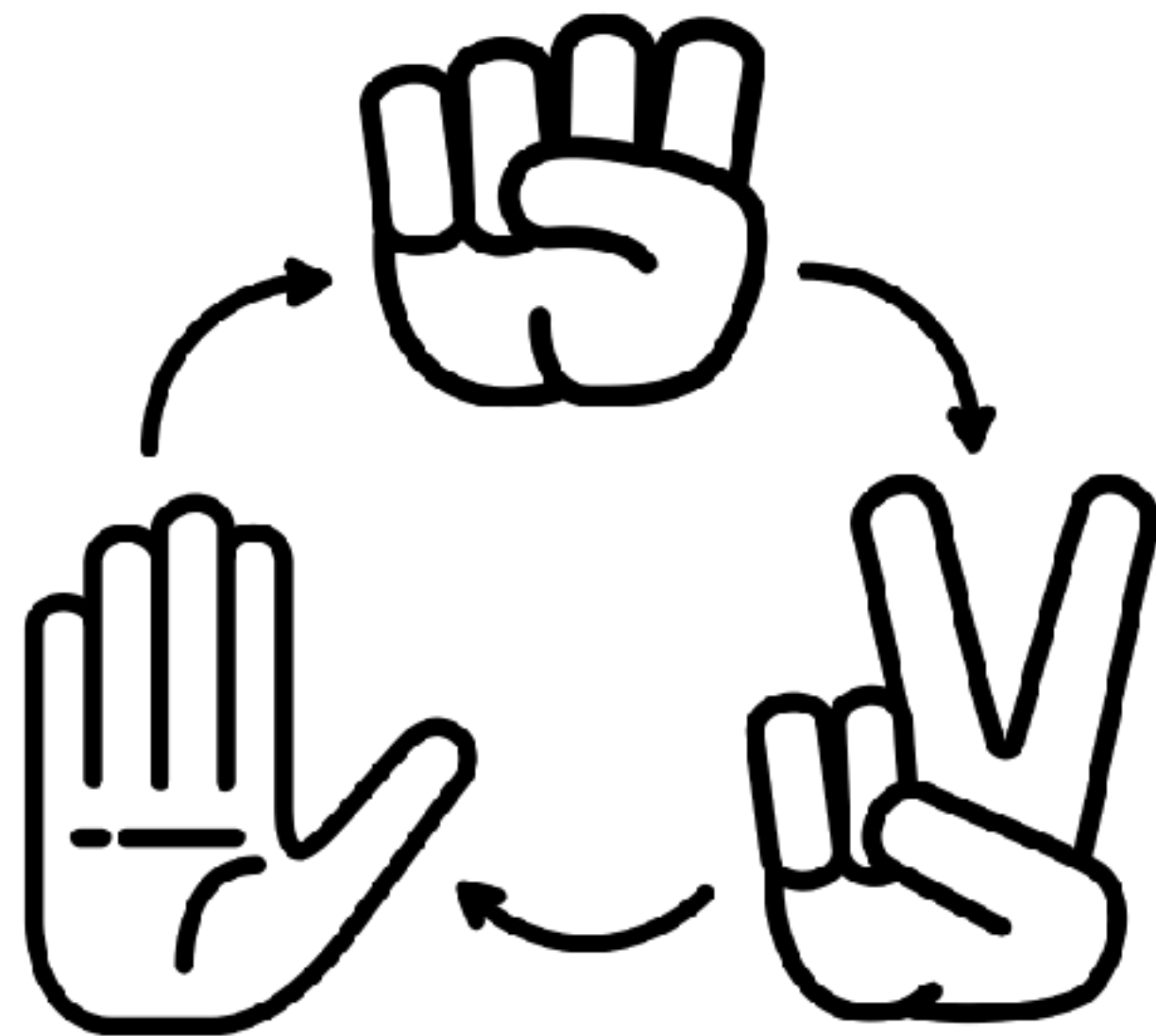


...but “mixed” Nash exist in finite spaces



(Bi)matrix game $\begin{cases} \min_{p \in \Delta} p^\top A q \\ \min_{q \in \Delta} p^\top B q \end{cases}$

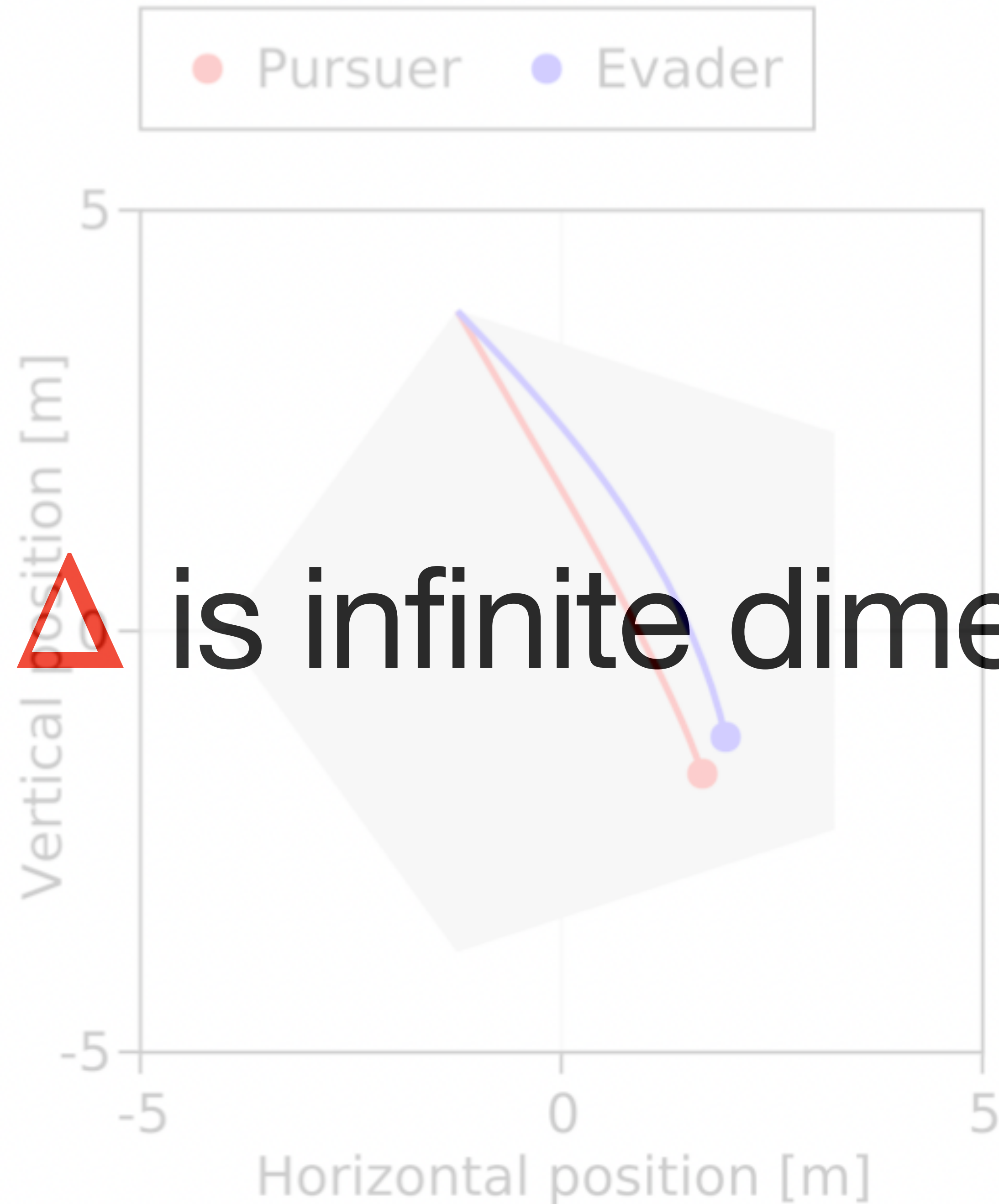
...but “mixed” Nash exist in **finite** spaces



(Bi)matrix game $\begin{cases} \min_{p \in \Delta} p^\top A q \\ \min_{q \in \Delta} p^\top B q \end{cases}$

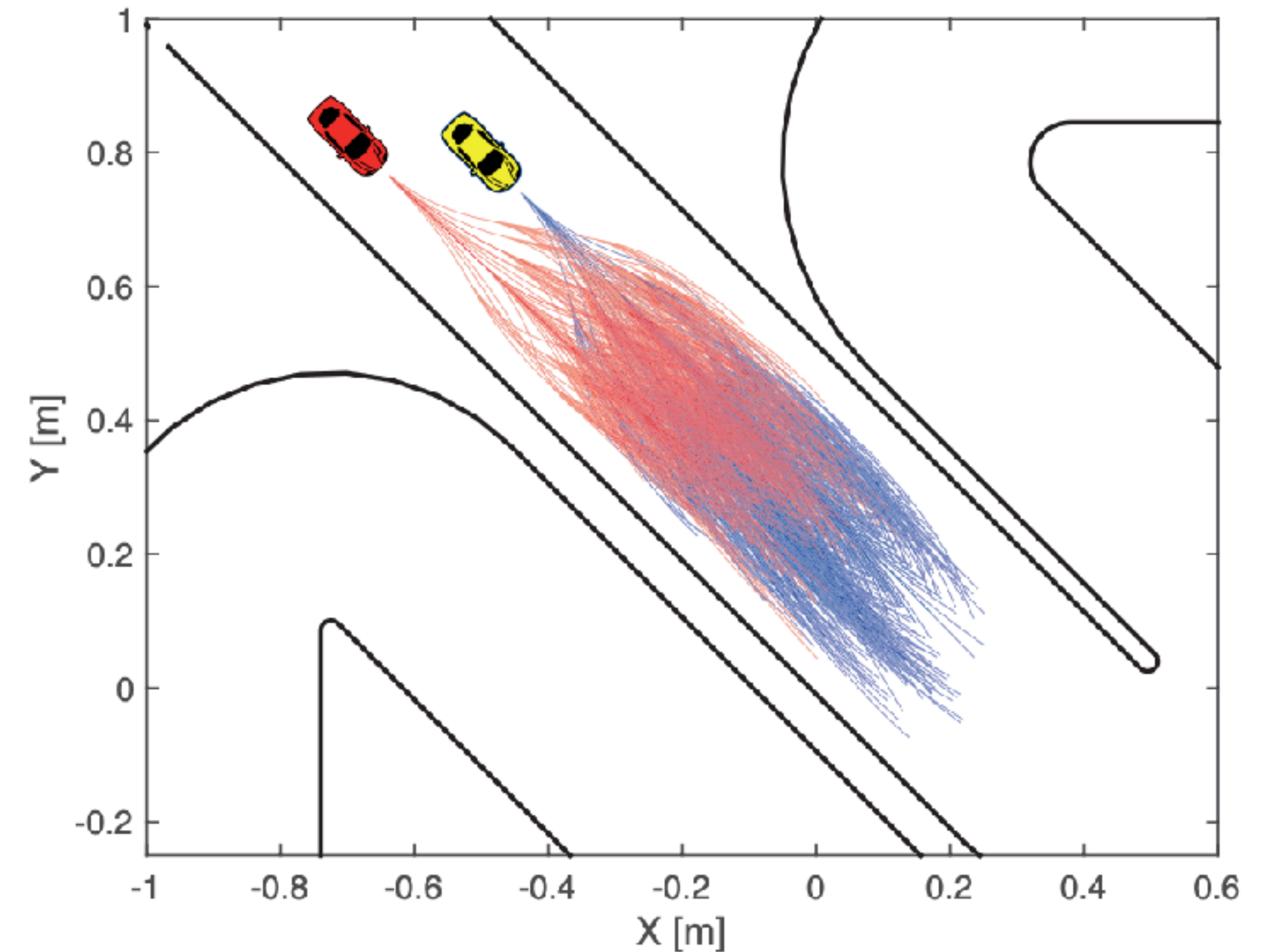


What if Δ is infinite dimensional?



Approaching the infinite case...

- Can construct the τ_j^i by dense sampling*
 - Not super efficient... big matrix game
- Observation 1: Sampling **more trajectories** yields quantifiably better performance#
- Observation 2: Mixed Nash solutions are almost always **sparse!**



*A. Liniger and J. Lygeros, "A Noncooperative Game Approach to Autonomous Racing," in *T-CST*, 2020.

#Bopardikar et. al., "Randomized Sampling for Large Zero-Sum Games," in *Automatica*, 2013.

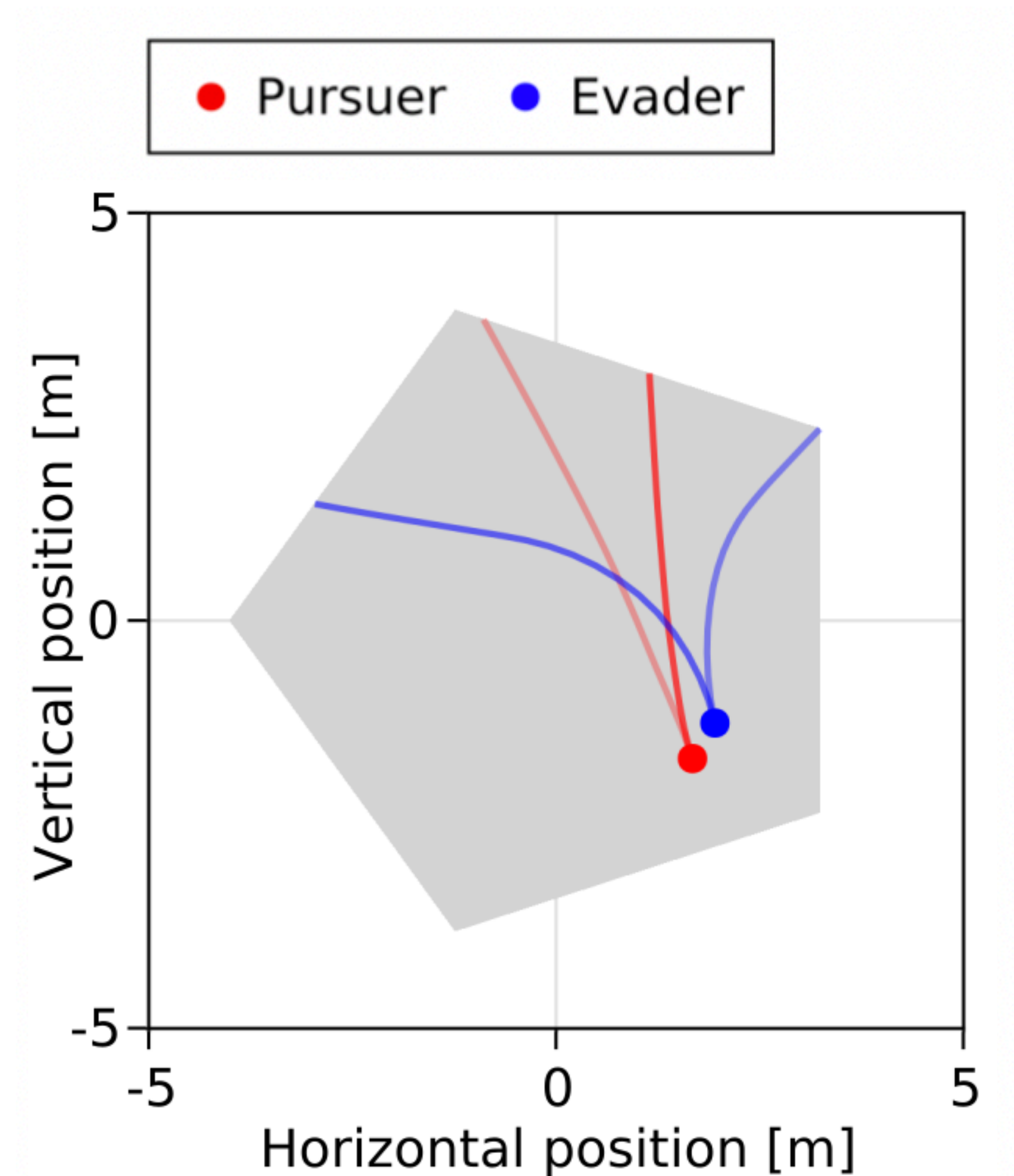
Idea: find **sparse** solutions directly
via a two-stage game

A two-stage formulation*

$$\left. \begin{array}{l} \min_{\{\tau_1^1, \dots, \tau_n^1\} \in \mathcal{T}_1^n} p^{*\top} A q^* \\ \max_{\{\tau_1^2, \dots, \tau_m^2\} \in \mathcal{T}_2^m} p^{*\top} A q^* \end{array} \right\} \text{Upper stage}$$

subject to (p^*, q^*) a Nash equilibrium

where $A_{ij} = f(\tau_i^1, \tau_j^2)$

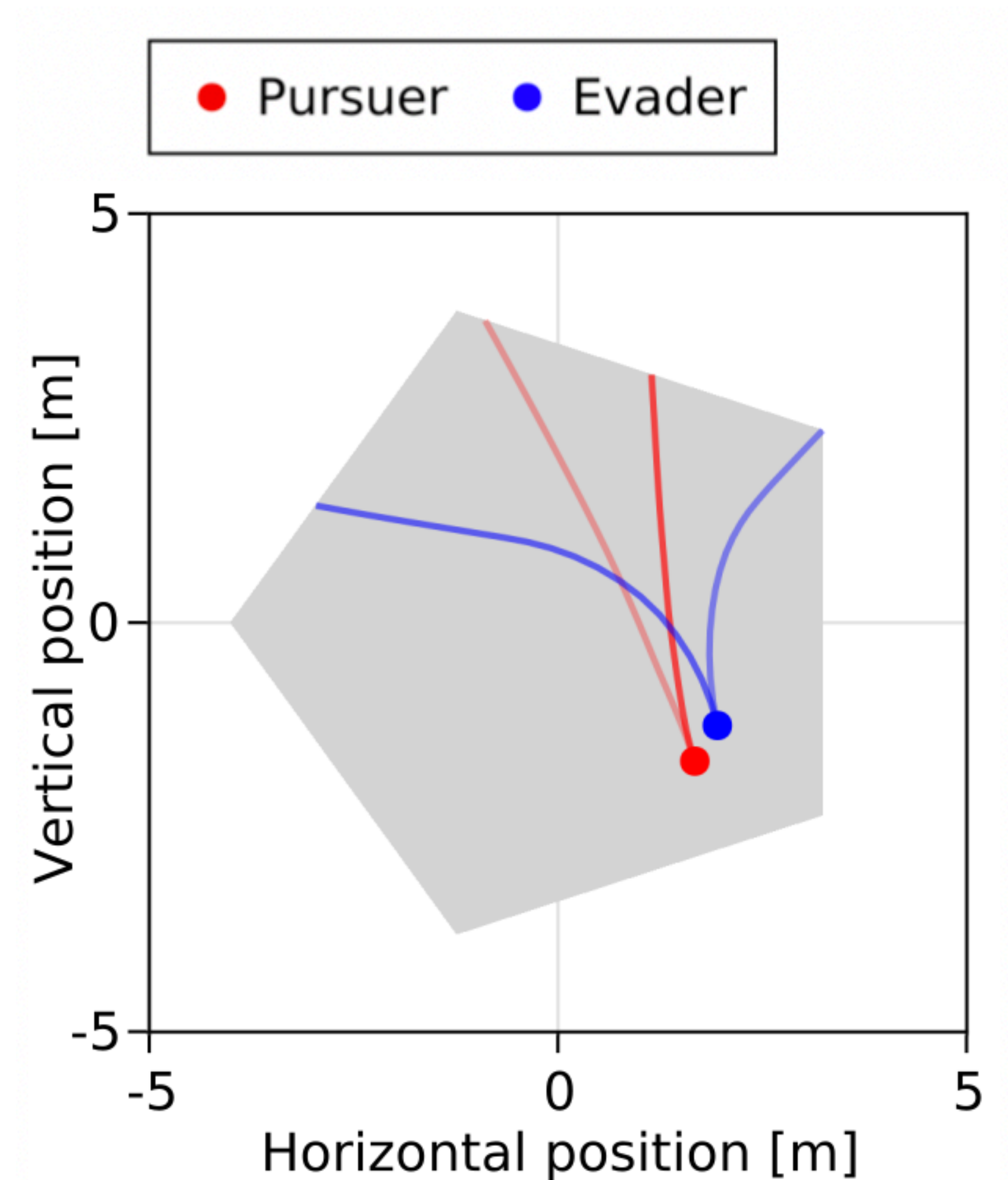


*L. Peters, **DFK**, et. al., "Learning Mixed Strategies in Trajectory Games" in RSS 2022.

A two-stage formulation*

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$$\left. \begin{array}{l} \text{subject to } p^* \in \arg \min_{p \in \Delta} p^\top A q \\ q^* \in \arg \max_{q \in \Delta} p^\top A q \end{array} \right\} \text{Lower stage}$$

where $A_{ij} = f(\tau_i^1, \tau_j^2)$



*L. Peters, **DFK**, et. al., "Learning Mixed Strategies in Trajectory Games" in RSS 2022.

A gradient-based solution method

$$\left. \begin{array}{l} \min_{\{\tau_1^1, \dots, \tau_n^1\} \in \mathcal{T}_1^n} p^{*\top} A q^* \\ \max_{\{\tau_1^2, \dots, \tau_m^2\} \in \mathcal{T}_2^m} p^{*\top} A q^* \end{array} \right\} \text{Upper stage}$$
$$\left. \begin{array}{l} \text{subject to } p^* \in \arg \min_{p \in \Delta} p^\top A q \\ q^* \in \arg \max_{q \in \Delta} p^\top A q \end{array} \right\} \text{Lower stage}$$

Simultaneous gradient descent/ascent on
 $\tau_{1:n}^1$ and $\tau_{1:m}^2 \dots$

...while chaining through the solution map
at the lower level

Implicit differentiation of NE

- Need to differentiate backward through (bi)matrix game solve
- Bimatrix game ~ linear complementarity program
- Can still differentiate solution map implicitly

$$(q_1, q_2) \text{ a NE for fixed } (A, B) \iff \begin{aligned} q_1 &\geq \mathbf{0} \perp Aq_2 \geq \mathbf{1} \\ q_2 &\geq \mathbf{0} \perp B^\top q_1 \geq \mathbf{1} \end{aligned}$$

strict complementarity \implies

$$\frac{\partial \bar{q}_2}{\partial \bar{A}} = \frac{\partial}{\partial \bar{A}} [\bar{A}^{-1} \mathbf{1}]$$

$\bar{A} \bar{q}_2 = \mathbf{1}$
 $\bar{B}^\top \bar{q}_1 = \mathbf{1}$
 $\hat{q}_2 = \mathbf{0}$
 $\hat{q}_1 = \mathbf{0}$

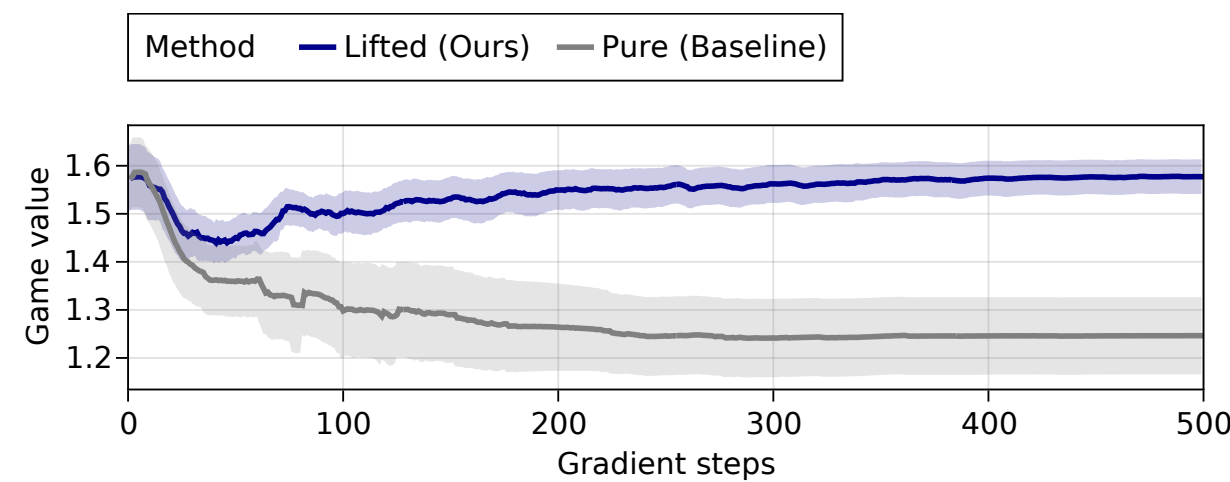
\implies

only active constraints

Be careful if only weak complementarity holds!

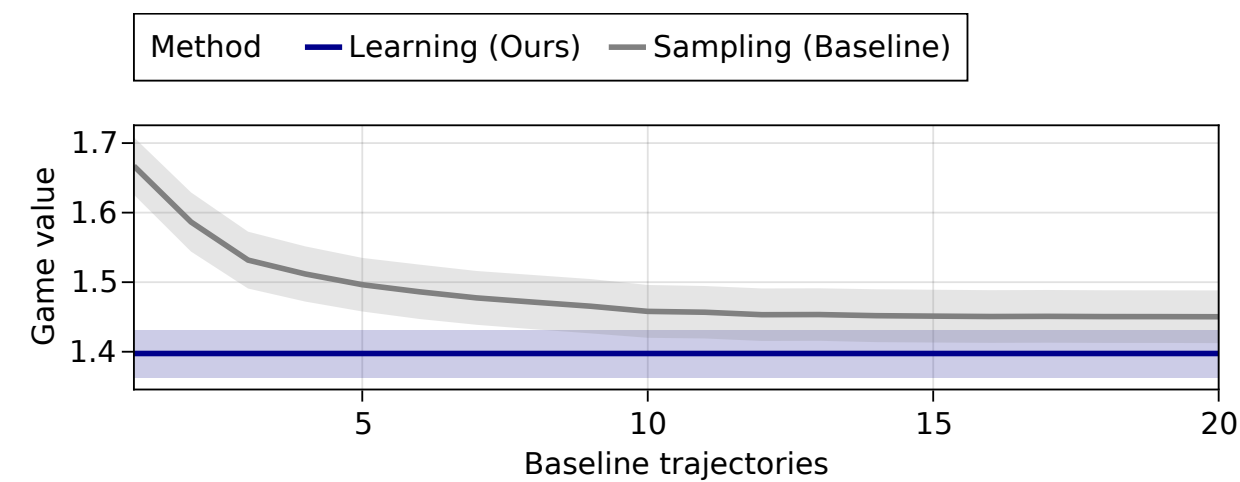
What happens?

Equilibrium characteristics



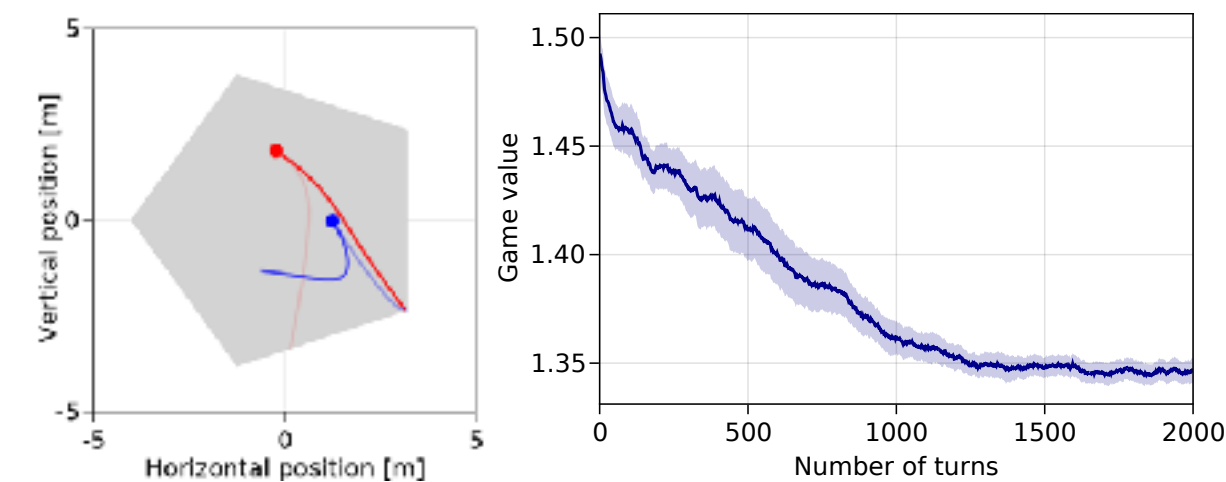
Lifting reliably converges to solutions which are NOT pure NE

Representative power



It's better to learn a few trajectory candidates than to sample many at random.

Online optimization in self-play



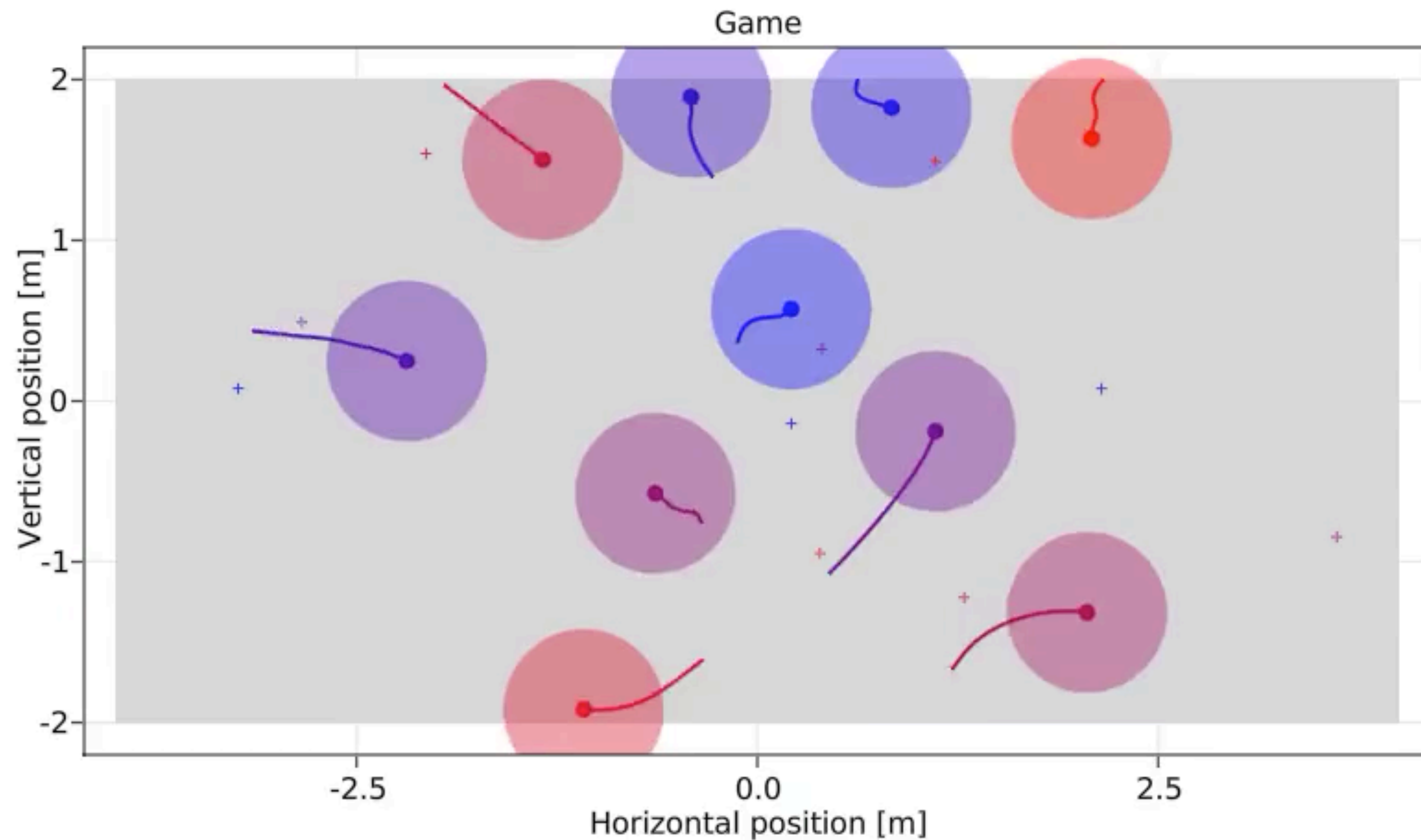
Lifted solvers can be pre-trained efficiently in self-play.

Meta-game: lifted vs. pure Nash

		Evader	
Pursuer	Lifted	Pure	
Lifted	1.360 ± 0.003	1.289 ± 0.005	
Pure	1.463 ± 0.004	0.903 ± 0.009	

Lifted strategies reliably outperform pure strategies.

Epilogue: the N-player case



- N-player setting is not a bimatrix game (see [TensorGames.jl](#))
 - 80ms to solve 6-player game with 3 actions each
- Parallelization + mixed-mode AD... pipeline scales very well for many players (single action):
 - 175ms per gradient step

Work with [Maximilian Schmidt](#)

Pure Nash solutions **do not always exist** in continuous trajectory games.
Mixed strategies depend upon **fixed** primitive trajectories.

Lifted strategies can be found efficiently via implicit differentiation,
and **outperform** pure counterparts.



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TU Delft



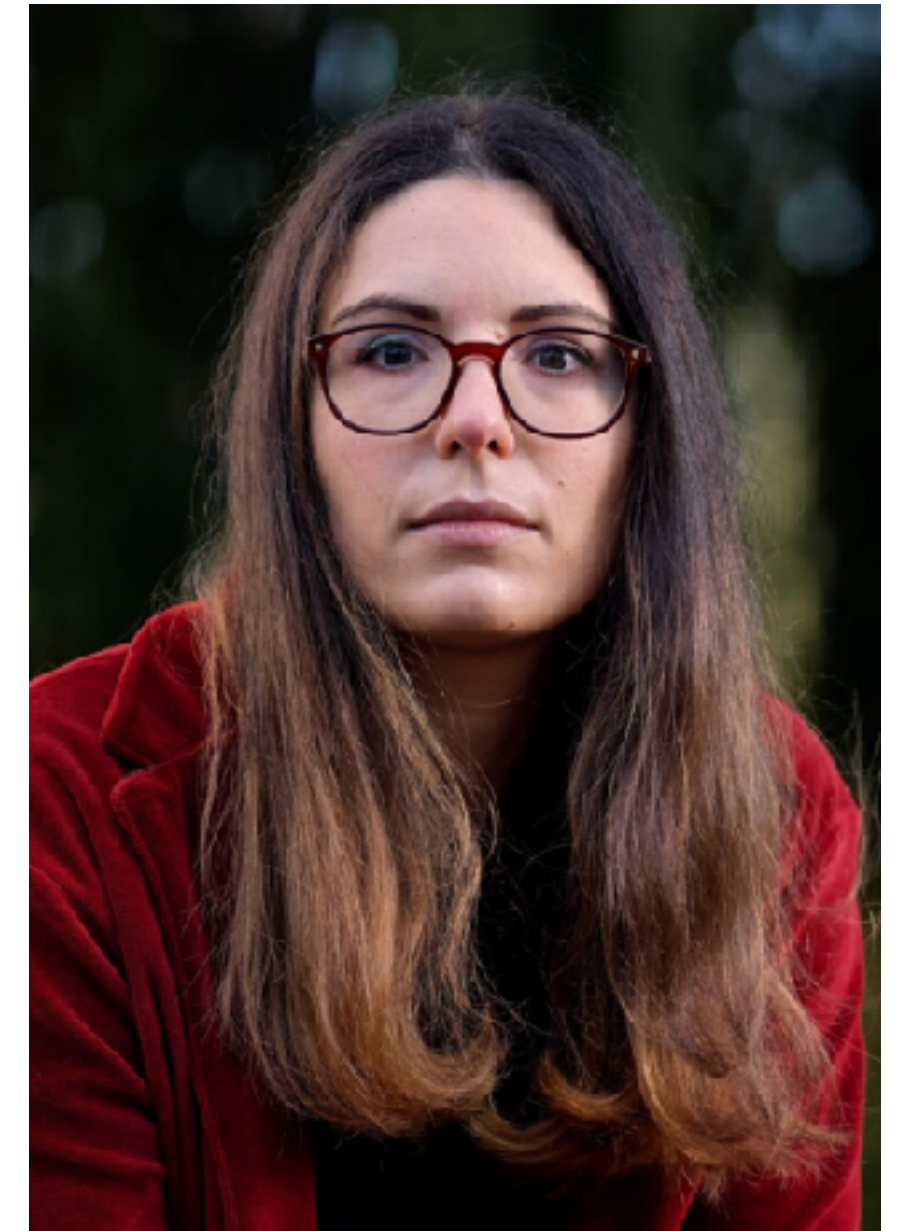
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Thank you! Questions?