

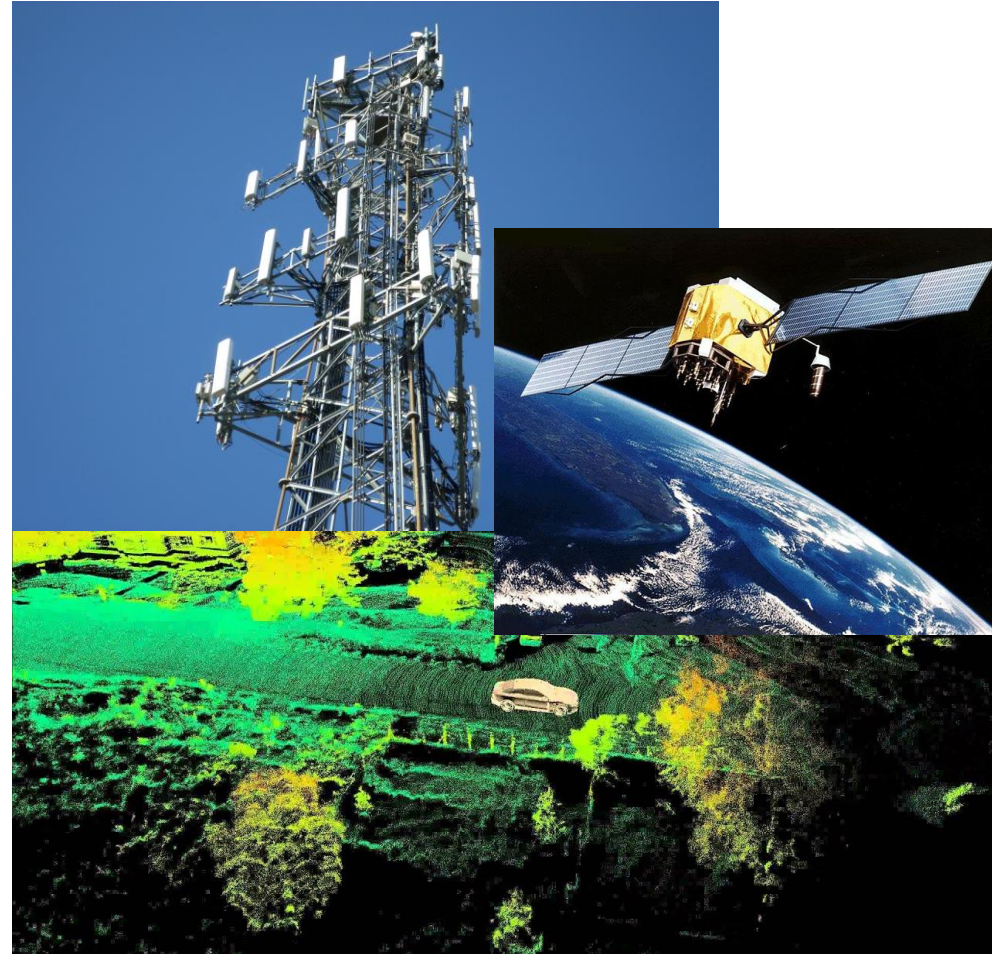
Path Planning in Environments with Intermittent State Feedback

S. C. Edwards, D. M. Le, D. P. Guralnik and W. E. Dixon, "A Topologically Inspired Path-Following Method With Intermittent State Feedback," in IEEE Robotics and Automation Letters, vol. 6, no. 3, pp. 4449-4456, July 2021, doi: 10.1109/LRA.2021.3067295.





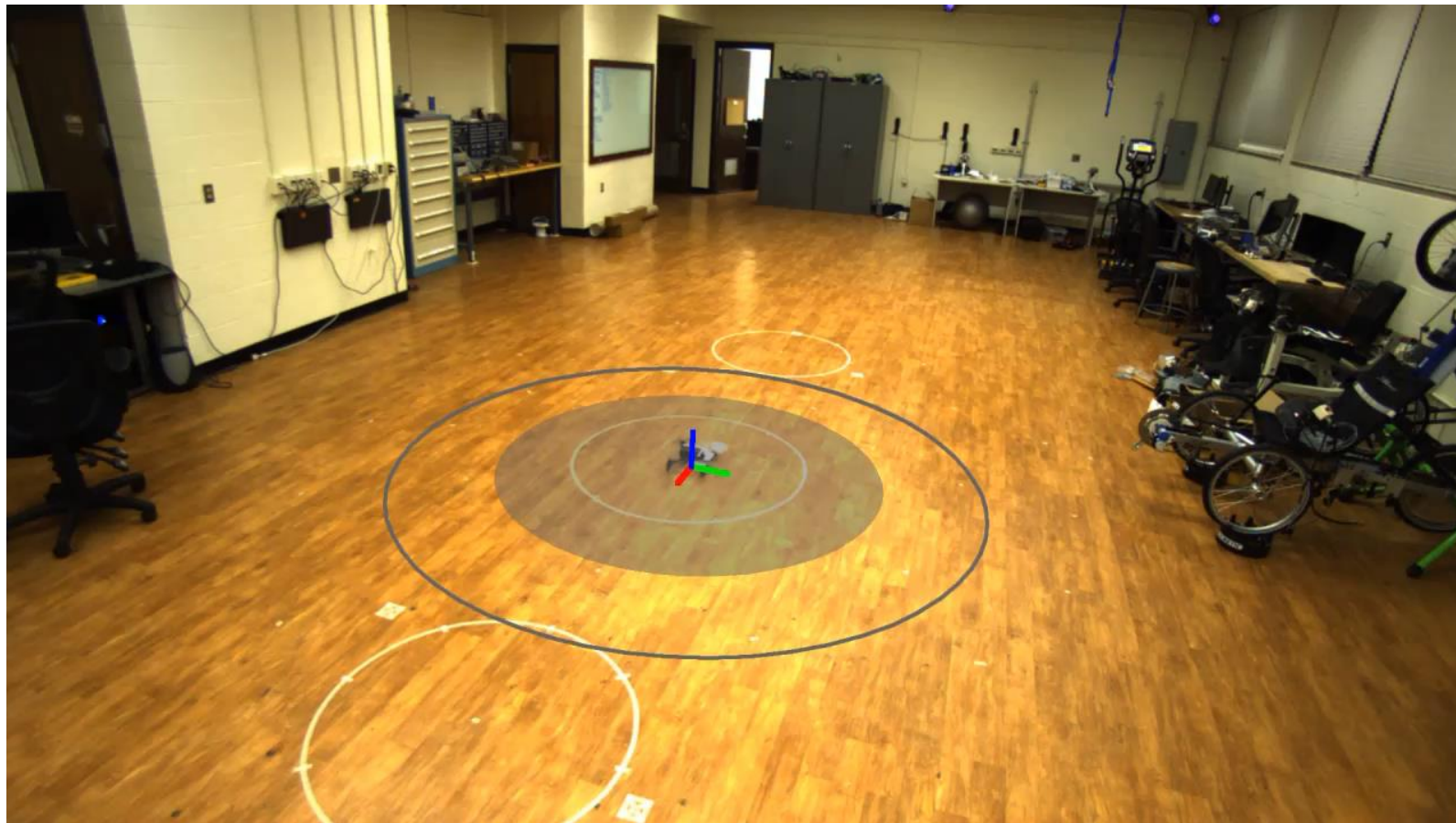
- Systems often operate in environments that include:
 - A2AD (Anti Access/Area Denial) Environments
 - Unknown Terrain
 - Complex Environments (Mountains, Extreme Weather, Marine, etc.)
- Limitations:
 - Sensing (GPS, Cameras, Lidar, etc.)
 - Global Communication





Problem Statement

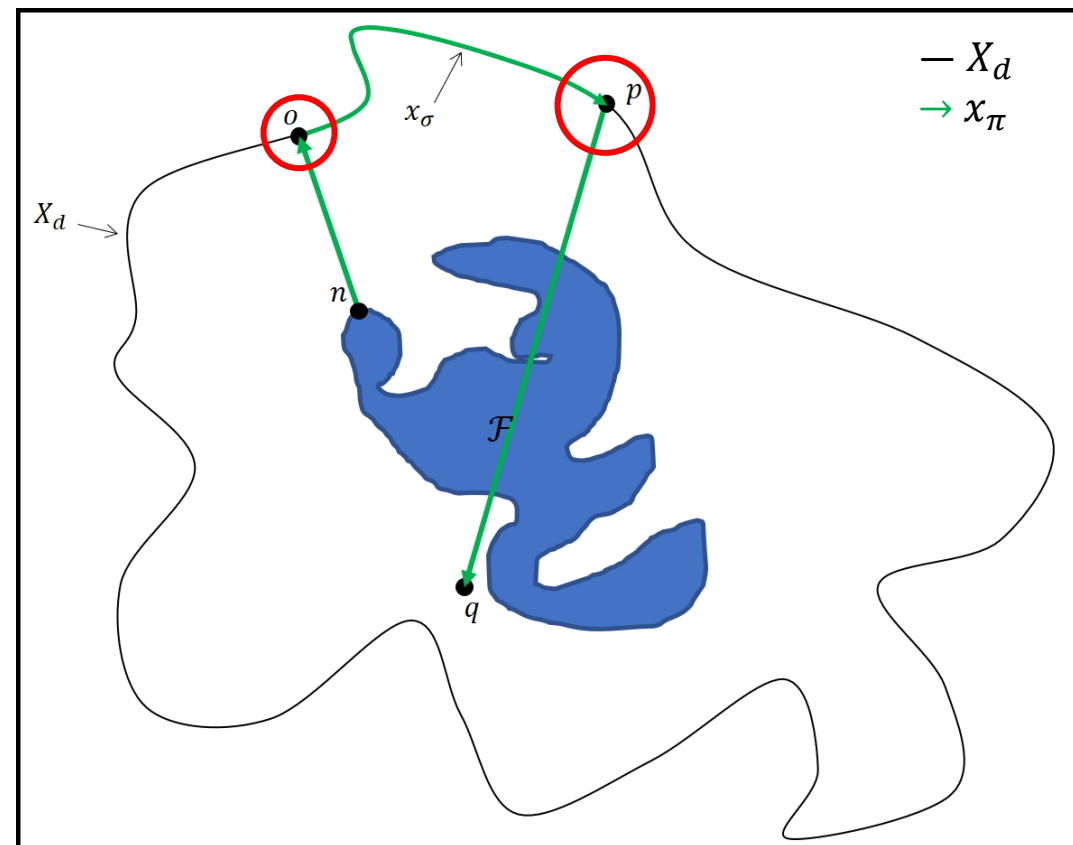
- Task: Agent is to follow path (X_d) where state feedback is denied
- Agent must switch between two tasks:
 - Following X_d
 - Obtaining state feedback to regulate tracking errors



H-Y. Chen, et. al., "A Switched Systems Approach to Path Following with Intermittent State Feedback", IEEE TRO, 2019.

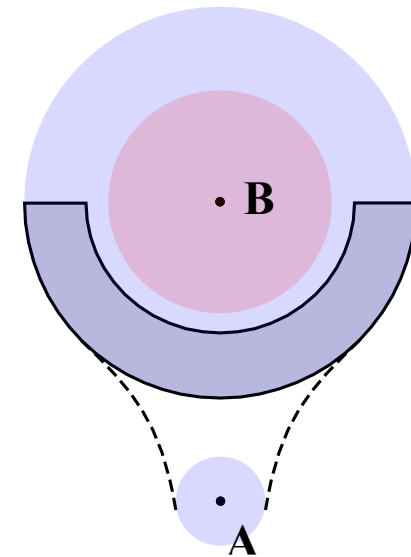
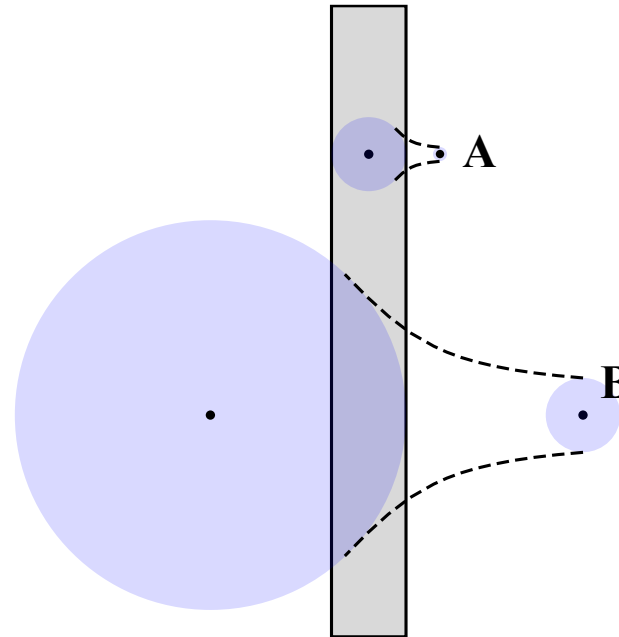
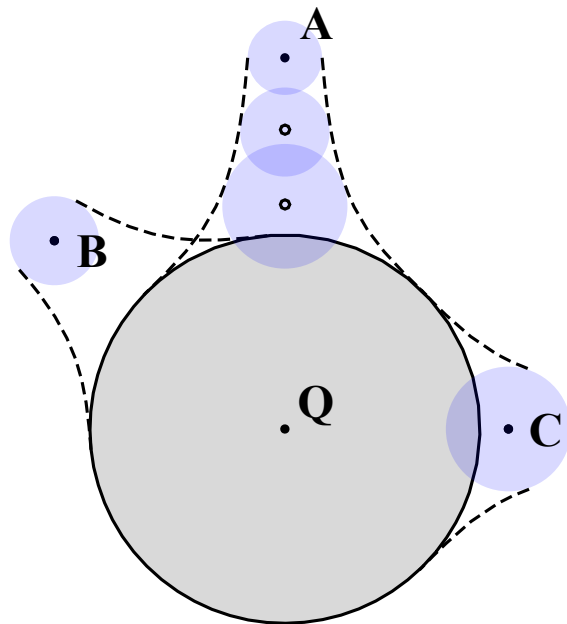


1. Depart from the feedback region to X_d .
2. Follow X_d until it is time to return to the feedback region.
3. Take a recommended return trajectory such to guarantee re-entry into the feedback region.

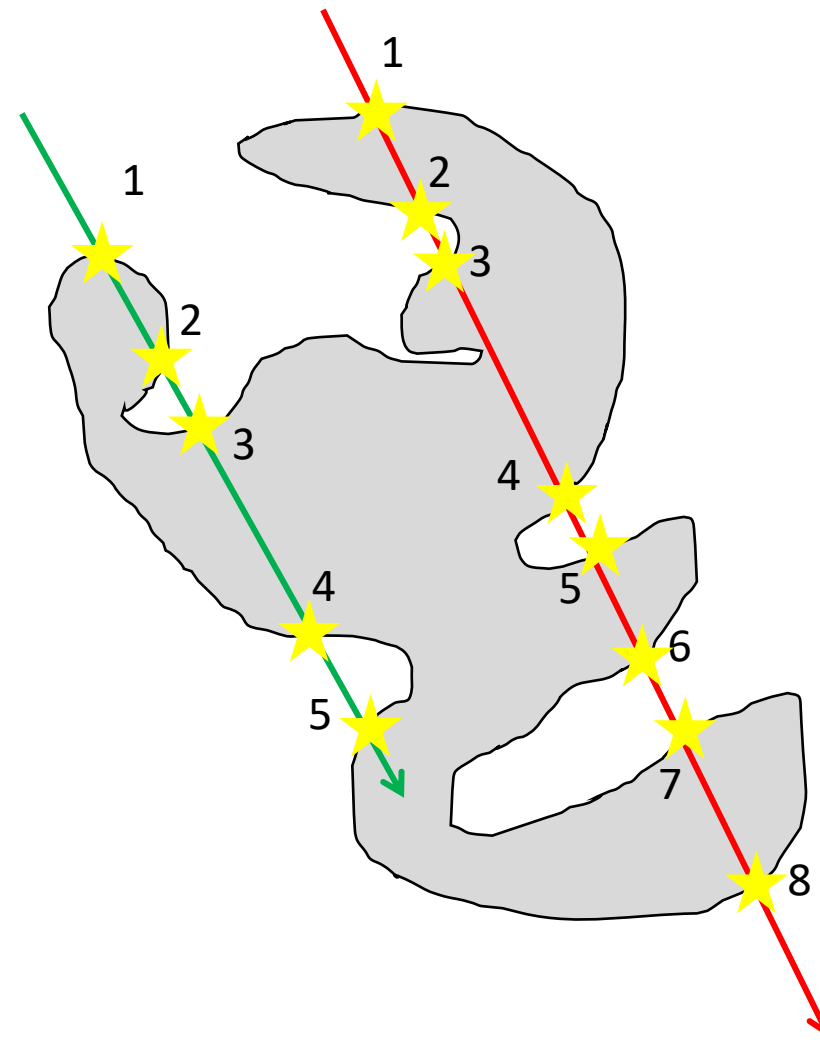




- Previous methods only considered circular feedback regions and circular paths to follow.
- This has issues...

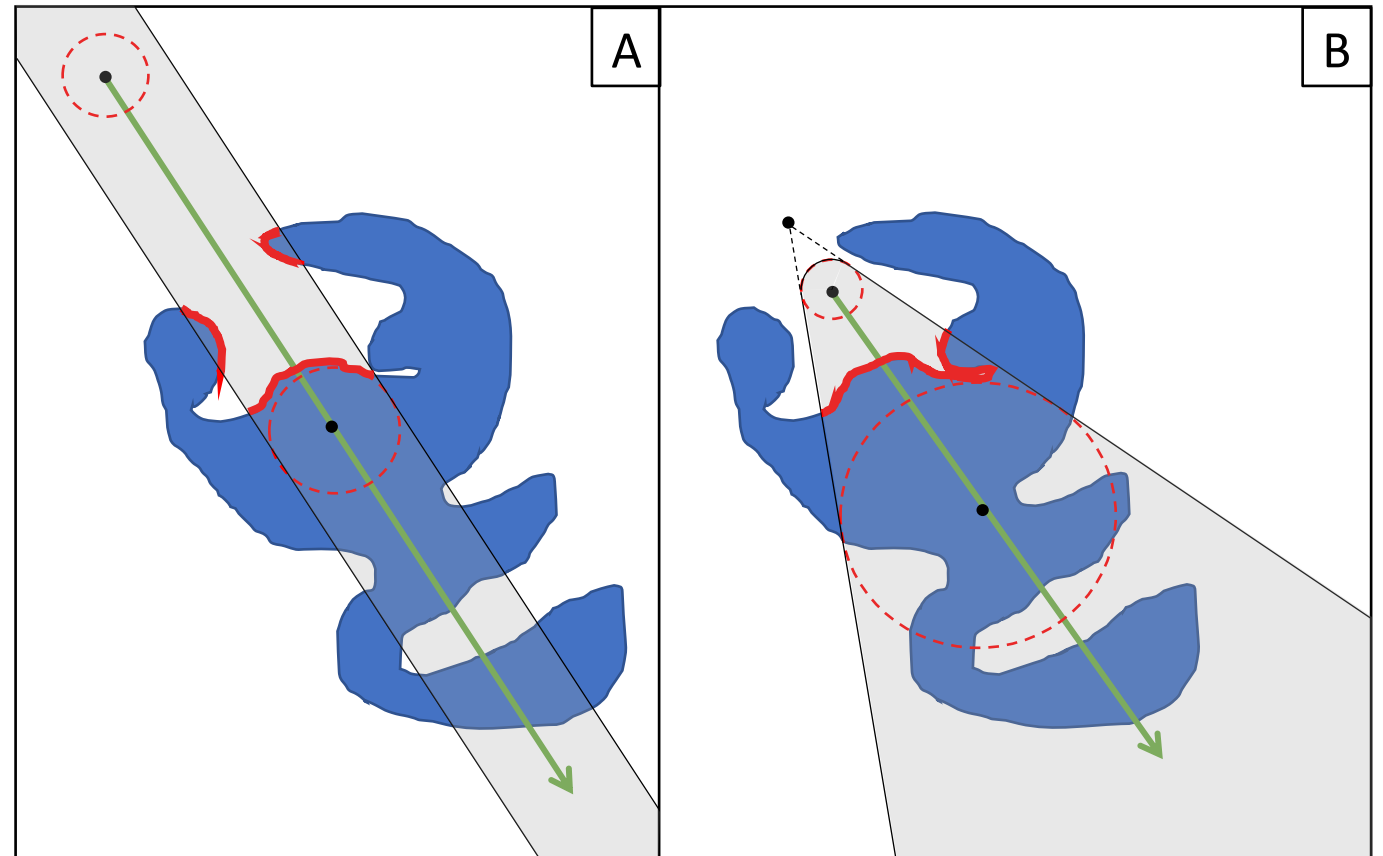


- Jordan Curve (C) separates plane into two regions
- A curve originating in the exterior with an odd number of intersections with C terminates in the interior
- An even number of intersections terminates in the exterior.



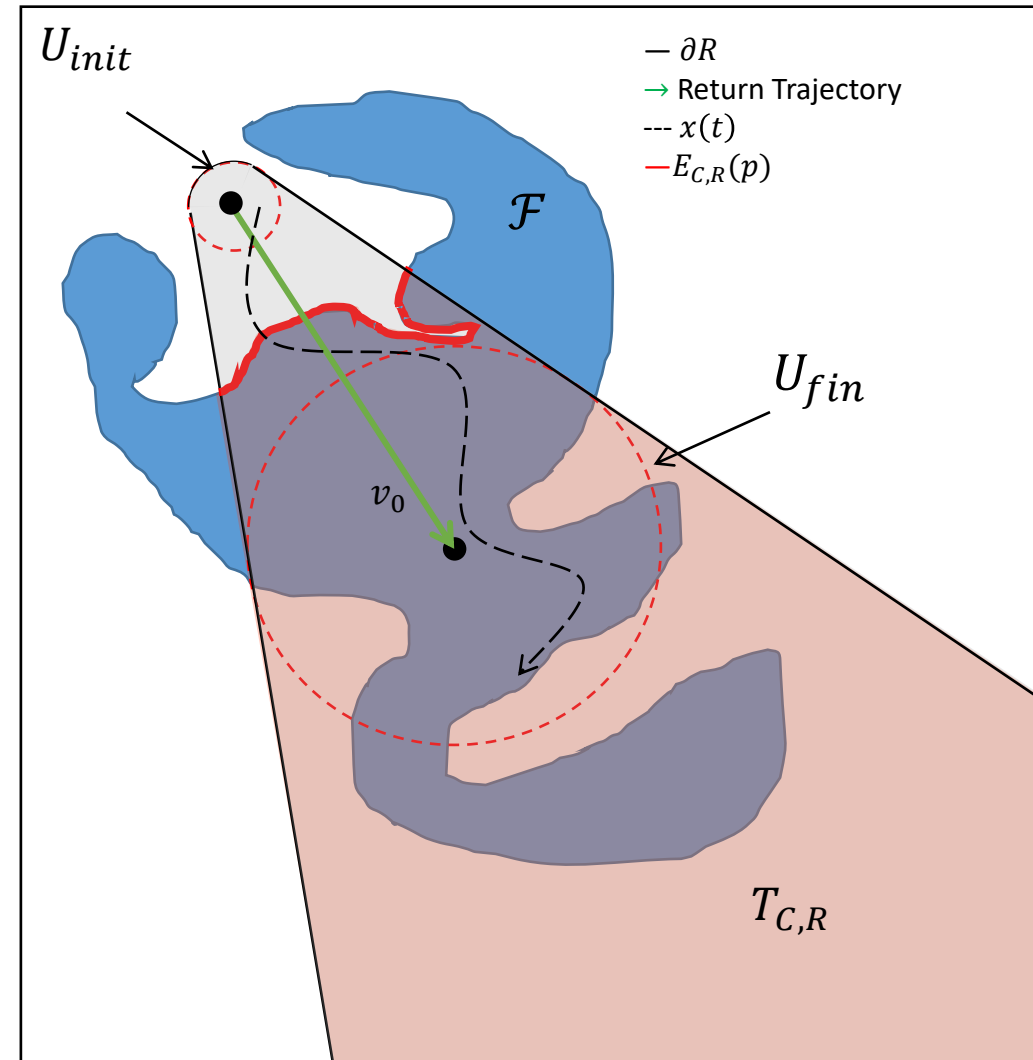


- Bounding regions are used to upper bound potential trajectories
- Required for separation of the initial position state and terminal state



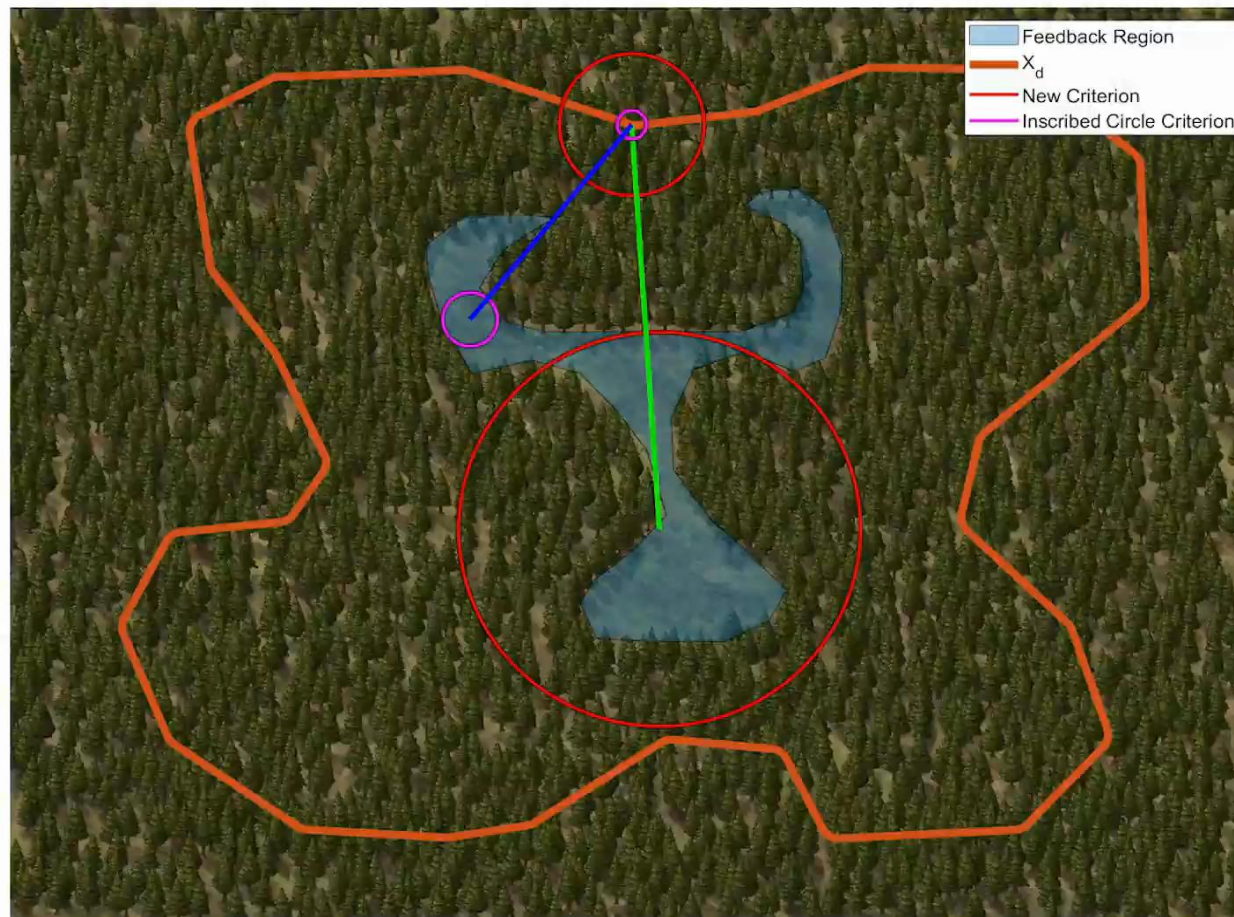


- Using Jordan Curves, “Target Regions” ($T_{C,R}$) are developed to guarantee an agent’s re-entry.
- The return target is the center of the ball inscribed in $T_{C,R}$.
- This generalizes to higher dimensions.



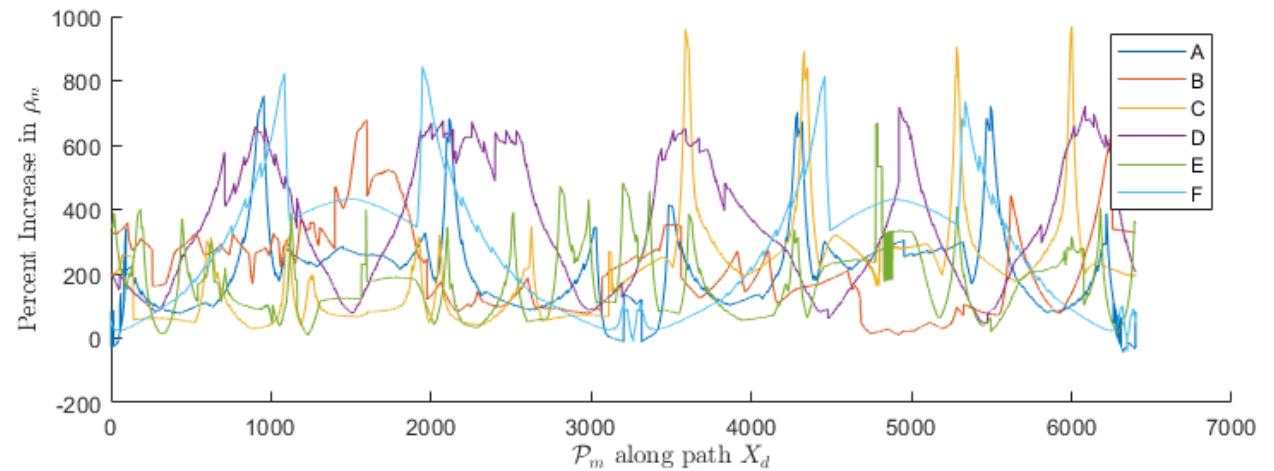
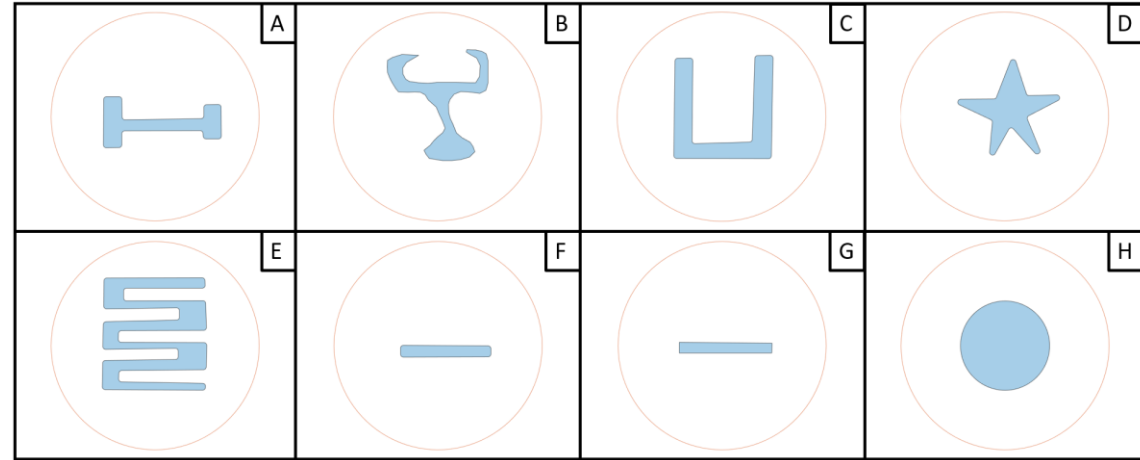


Maximum Allowed Uncertainty Regions (MAURs)





- With a total of 46,855 data points, the mean-average increase in the MAUR was found to be 233%.
- The largest increase was found to be 969%. This was observed in the “horseshoe” geometry (C).





MAUR Computation Methods

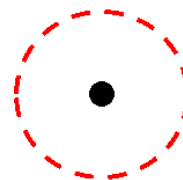
- Current methods use a brute force approach
- New methods aim to reduce computation time from hours to fractions of a second
- Add other animation of old method

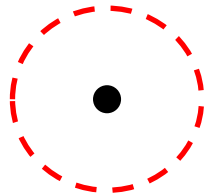


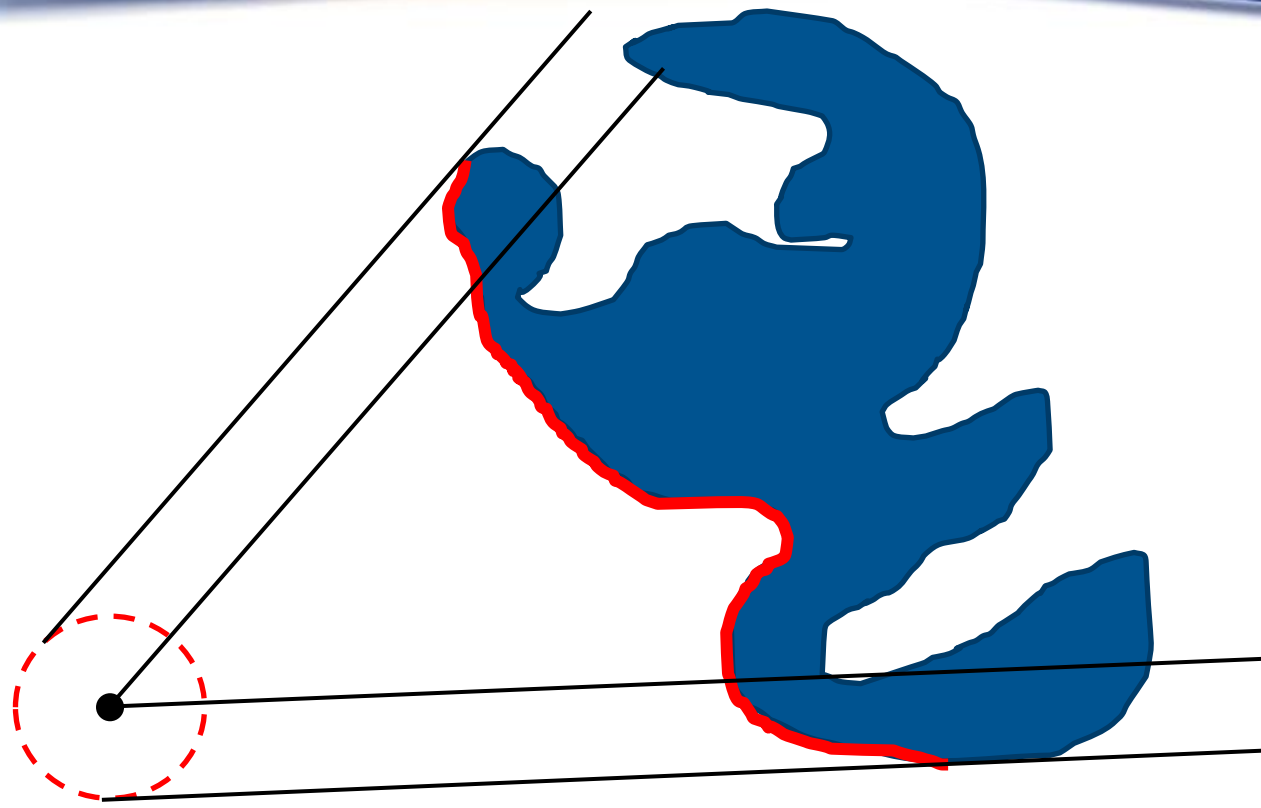


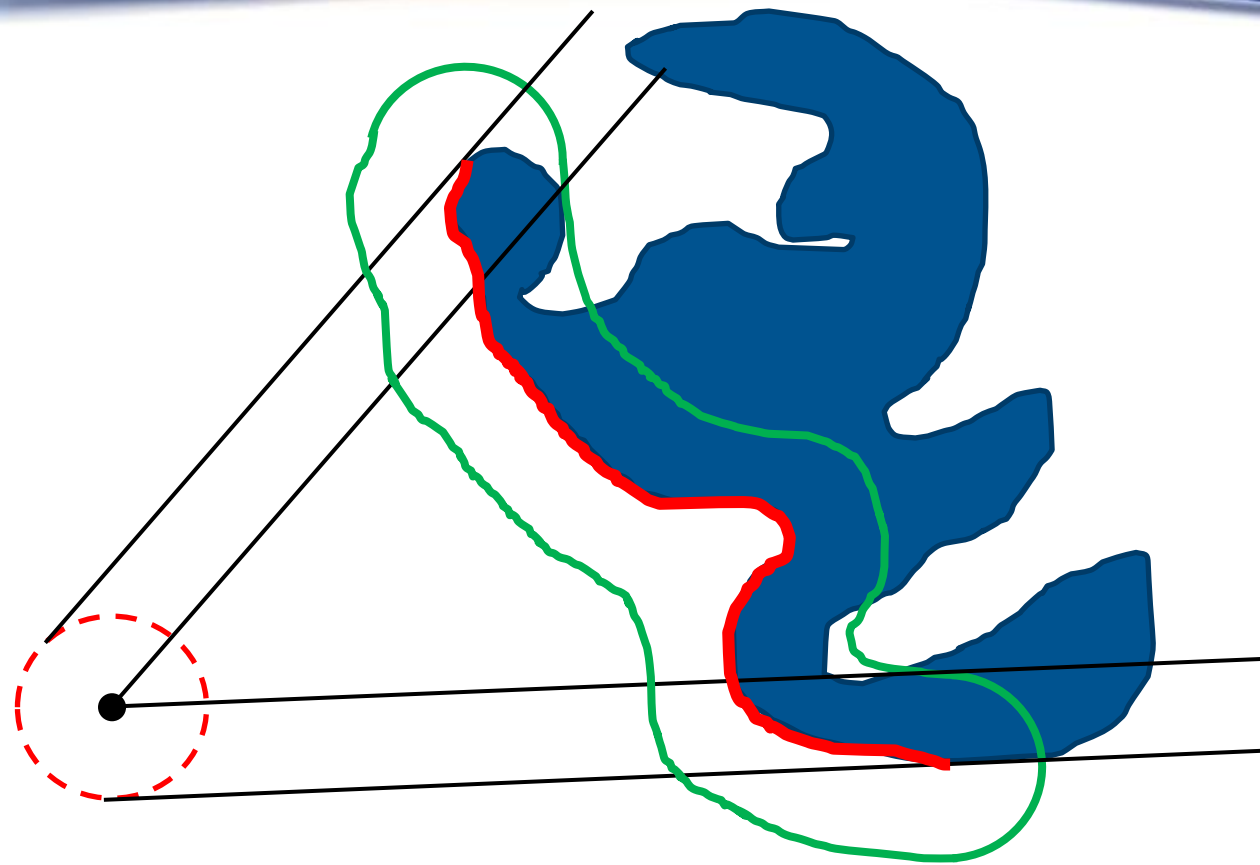
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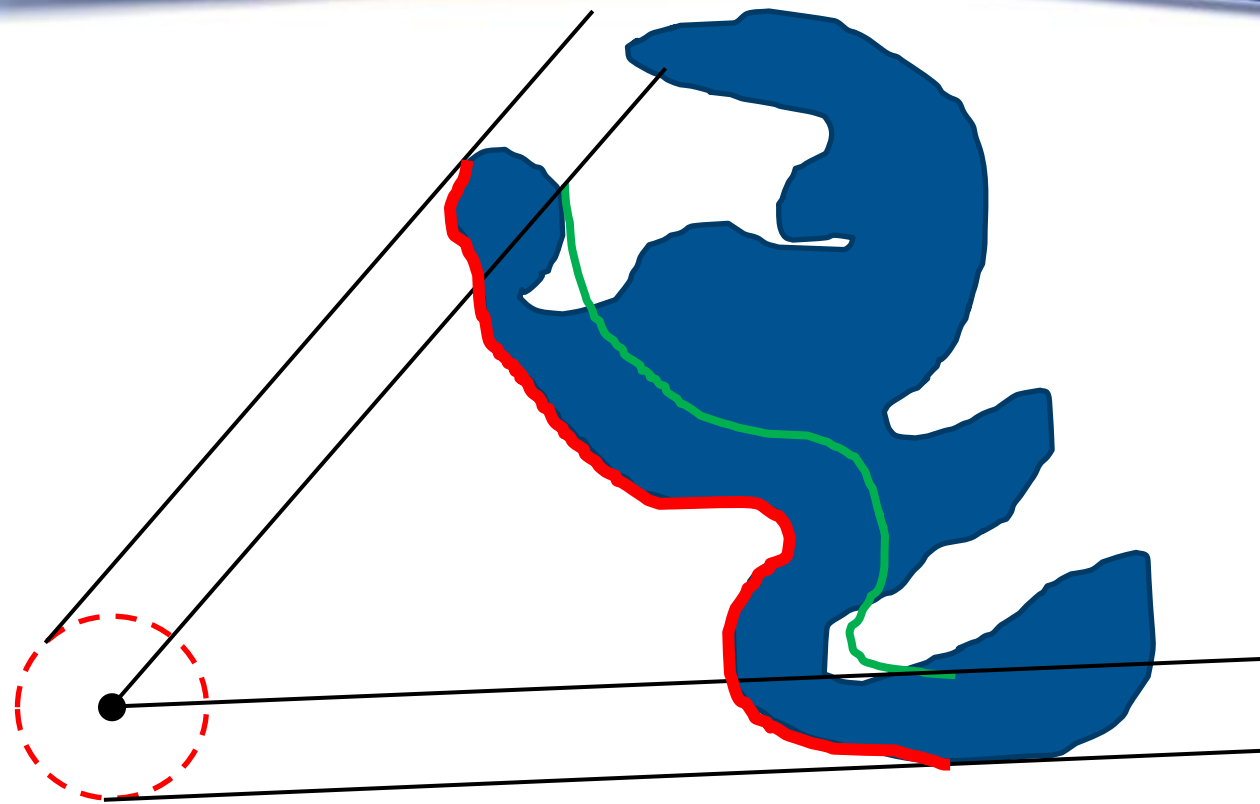
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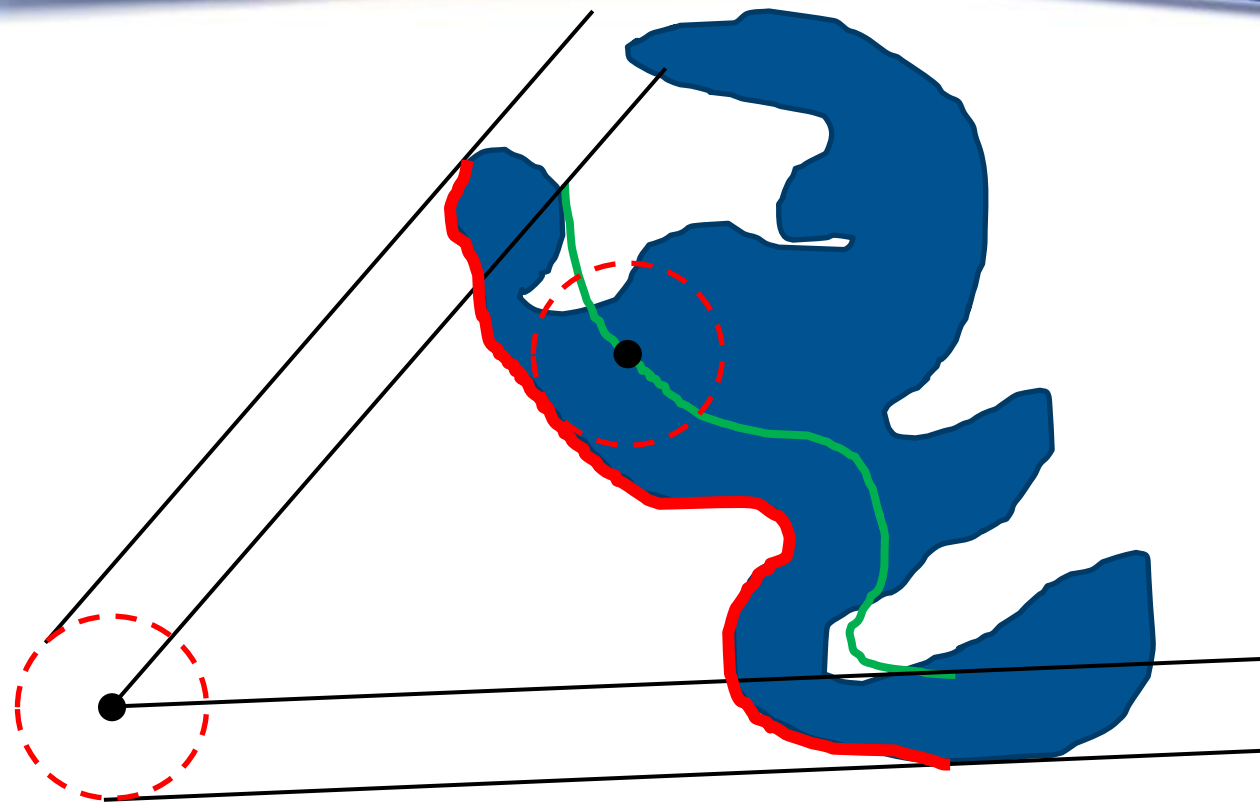












Topological Methods for Guaranteeing Transitions in Switched Systems

Dan P. Guralnik

University of Florida/NCR Lab

Sage C. Edwards

Warren E. Dixon

April 30, 2021

Informal Challenge: *Given a controlled hybrid system with uncertain state, how to guarantee transitions into a desired operational mode?*

~> *crucial for ensuring high-level plan execution*

~> *complicated geometry of transition boundaries*

~> *must be addressed with minimal specific geometric insight*



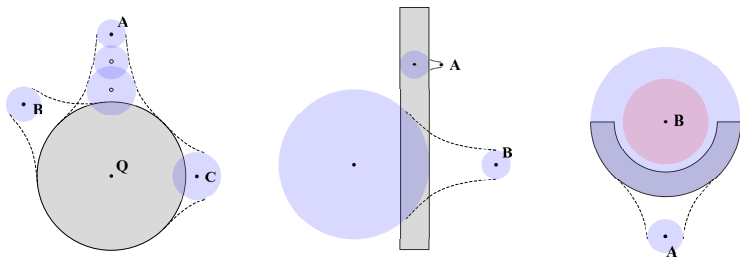
Motivation

Informal Challenge: *Given a controlled hybrid system with uncertain state, how to guarantee transitions into a desired operational mode?*

~> *crucial for ensuring high-level plan execution*

~> *complicated geometry of transition boundaries*

~> *must be addressed with minimal specific geometric insight*



Guaranteed Transitions: Lack of symmetry makes the problem less intuitive...

~> *It's not about thickness of the target domain*

~> *It's about the transition boundary separating the error cone*



Formal Problem Statement

A setup, successfully generalizing our work in [1]:

- ▶ *Incomplete information on state-environment interaction*: agent dynamics over $x \in \mathcal{F}^0$, $\mathcal{F} \subset \mathbb{R}^n$ closed, given by

$$\|\dot{x} - f(x, u)\| \leq \bar{d}. \quad (1)$$



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- ▶ A collection \mathcal{A} of curves $\hat{x} : [0, \infty) \rightarrow \mathbb{R}^n$, called *admissible plans*,

$$\mathcal{A}_p \triangleq \{\hat{x} \in \mathcal{A} : \hat{x}(0) = p\}, \quad (2)$$

\rightsquigarrow Usually, expect to have $\mathcal{A}_p = p + \mathcal{A}_0 = \{p + \hat{x} : \hat{x} \in \mathcal{A}_0\}$

\rightsquigarrow In Sage's problem $\hat{x}(t) \triangleq p + t\mathbf{v}$, with $\|\mathbf{v}\| = v_0$ a known constant



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\rightsquigarrow In Sage's problem $\hat{x}(t) \triangleq p + t\mathbf{v}$, with $\|\mathbf{v}\| = v_0$ a known constant

- ▶ Each plan is *executed* by applying a control $u(t)$, $t \geq 0$, guaranteeing

$$\|x(t) - \hat{x}(t)\| \leq \varrho(t), \quad t \geq 0 \quad (3)$$

where $\varrho : [0, \infty) \rightarrow [0, \infty)$ is a known *error-bounding function* (EBF).

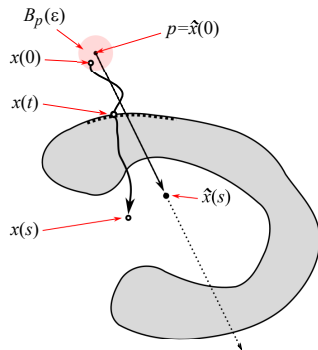
\rightsquigarrow There can be many EBFs, e.g., depending on $p = \hat{x}(0)$



Formal Problem Statement

Definition. A plan $\hat{x} \in \mathcal{A}_p$ provides an ε -guaranteed transition into \mathcal{F} , if

$$(\forall p^* \in B_p(\varepsilon)) (\exists t \geq 0) (x(0) = p^* \rightarrow x(t) \in \mathcal{F}). \quad (4)$$



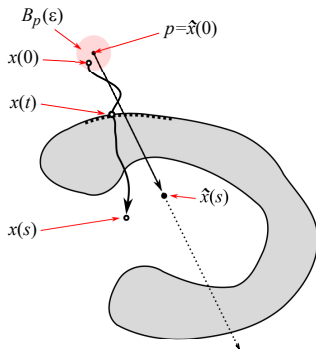


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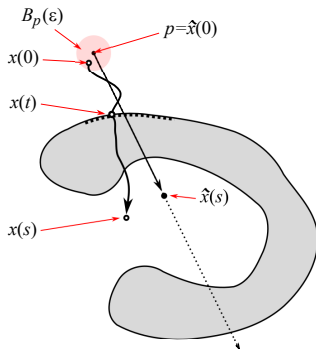


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- ▶ $p \in \mathbb{R}^n$ is ε -feasible, if $\mathcal{G}_p(\varepsilon) \neq \emptyset$.
- ▶ $p \in \mathbb{R}^n$ is *feasible* if it is ε -feasible for some $\varepsilon > 0$.
- ▶ **WHAT ACTUALLY GUARANTEES A PLAN?**



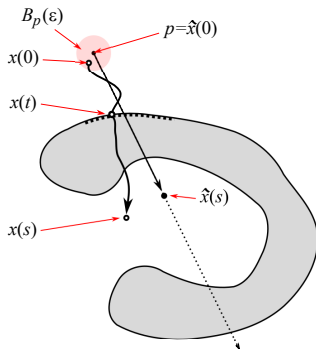


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- ▶ **WHAT ACTUALLY GUARANTEES A PLAN?**
- ▶ **Recall:** Let $U \subset \mathbb{R}^n$ be an open domain and let $p, q \in U$. A closed set $K \subset \mathbb{R}^n$ *separates U between p and q* , if p, q lie in distinct components of $U \setminus K$.
- ▶ $\partial \mathcal{F}$ separates \mathbb{R}^n between $p \in \mathcal{F}^0$ and any $q \in \text{int}(\mathcal{F})$.





Formal Problem Statement

Definition. The *maximum allowed uncertainty radius* (MAUR) at p :

$$\mathbf{M}_{\mathcal{F}}^*(p) \triangleq \begin{cases} \sup(\varepsilon > 0: \mathcal{G}_p(\varepsilon) \neq \emptyset), & \text{if } p \text{ is feasible,} \\ 0 & \text{otherwise,} \end{cases} \quad (5)$$

and denote $\mathcal{M}(\mathcal{F}) \triangleq \{f: \mathbb{R}^n \rightarrow \mathbb{R}_{\geq 0} : f \leq \mathbf{M}_{\mathcal{F}}^*\}$, the *MAUR bounds*.

Problem 1. Find explicit constructions of $f \in \mathcal{M}(\mathcal{F})$.

↔ The inscribed ball criterion (IBC) leads to one such construction

Problem 2. Given $f \in \mathcal{M}(\mathcal{F})$ and $p \in \mathbb{R}^n$, compute $f(p)$ and $\hat{x} \in \mathcal{G}_p(f(p))$.

↔ Expect: the better f approximates $\mathbf{M}_{\mathcal{F}}^$, the more complex this problem will be*

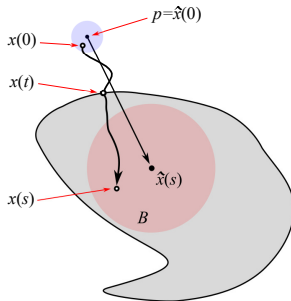


Topological Transition Guarantee (TTG)

► An *IBC certificate at p* is an EBF ϱ with

► $B_{\hat{x}(s)}(\varrho(s)) \subseteq \mathcal{F}$,

for some $\hat{x} \in \mathcal{A}_p$, $s \geq 0$.





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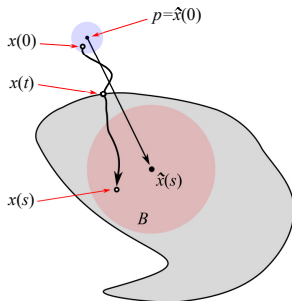
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- ▶ $M_{\mathcal{F}}^{IBC}(p) \triangleq \sup \left(\varrho(0) : \begin{array}{l} \varrho \text{ is an IBC} \\ \text{certificate at } p \end{array} \right)$

and $M_{\mathcal{F}}^{IBC}(p) \triangleq 0$ otherwise.





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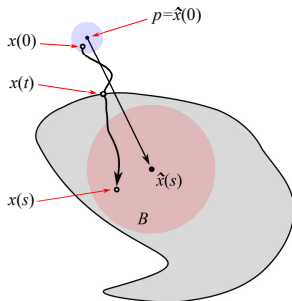
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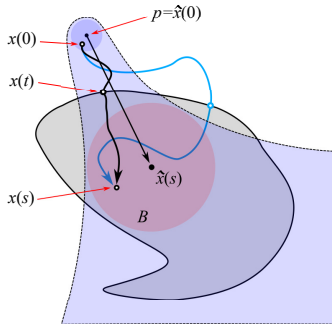
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- ▶ Clearly, $M_{\mathcal{F}}^{IBC} \leq M_{\mathcal{F}}^*$.
- ▶ The argument takes no account of the error cone!

↔ Information about $x(t)$ is thrown away



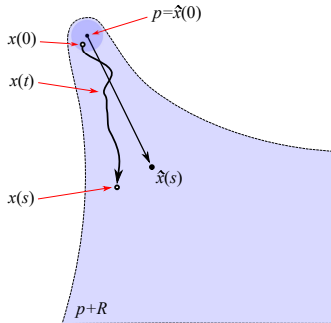


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► Let \mathcal{U} be a collection of open domains such that

- $\mathbb{R}^n \in \mathcal{U}$,
- $0 \in R$ for all $R \in \mathcal{U}$.

↪ In [1], \mathcal{U} consists of all strips centered at $0 \in \mathbb{R}^2$
↪ \mathcal{U} expresses a bound on the growth of error cones





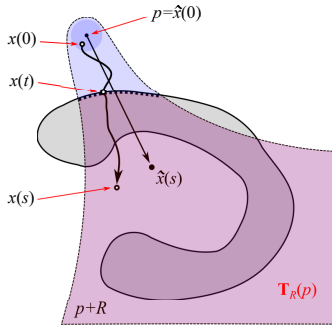
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- ▶ Given $R \in \mathcal{U}$, $p \in \mathbb{R}^n$ the **target at p** , $\mathbf{T}_R(p)$, is defined as the set of $q \in \mathbb{R}^n$ separated from p by $\partial\mathcal{F}$ in $p + R$.





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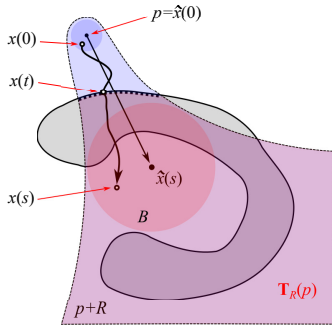
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- ▶ A **\mathcal{U} -certificate at p** is an EBF ϱ , with:

- ▶ $B \triangleq B_{\hat{x}(s)}(\varrho(s)) \subseteq \mathbf{T}_R(p)$, and
- ▶ $\hat{x}(t) \in p + R$ for all $0 \leq t \leq s$.

for some $\hat{x} \in \mathcal{A}_p$, $s \geq 0$ and $R \in \mathcal{U}$.



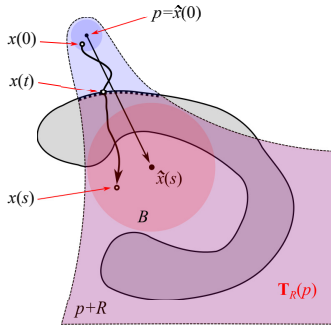


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► We can then define the \mathcal{U} -MAUR:

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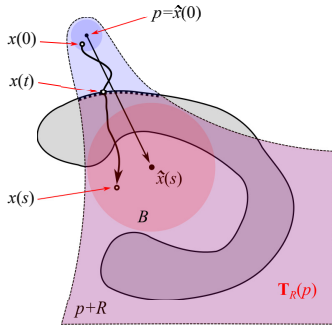
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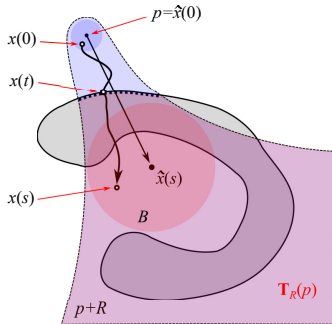
$$\text{► } M_{\mathcal{F}}^{IBC} \leq M_{\mathcal{F}}^{\mathcal{U}} \leq M_{\mathcal{F}}^*.$$





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- ▶ Since $\mathbf{T}_{\mathbb{R}^n}(p) = \mathcal{F}$, we have
 - ▶ $M_{\mathcal{F}}^{IBC} \leq M_{\mathcal{F}}^{\mathcal{U}} \leq M_{\mathcal{F}}^*$.
- ▶ HENCE A FOCUS ON COMPUTABILITY OF LOWER BOUNDS $f \leq M_{\mathcal{F}}^{\mathcal{U}}$, OBTAINED BY
 - ▶ restricting \mathcal{U} : strips, cylinders, cones
 - ▶ restricting ϱ : specific EBFs
 - ▶ restricting $\partial\mathcal{F}$: collared spheres





Towards Computable TTG MAURs

- ▶ *Jordan Curves* [2] provide a “separation standard” in the plane.
 - ▶ $\partial\mathcal{F}$ can be any continuous simple closed curve (SCC)
 - ▶ $\partial\mathcal{F}$ can be lower-approximated by polygonal SCCs



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- ▶ For simple parametric \mathcal{U} (strips, cones), lower bounds on $M_{\mathcal{F}}^{\mathcal{U}}(p)$ are...
 - ▶ Piecewise regular for regular increasing ϱ
 - ▶ Solutions to a constrained optimization problem
 - ~> *Combinatorial/Topological aspects to be resolved*
 - ▶ Computations are parallelizable.
 - ~> *Target balls picked via edge-by-edge optimization*



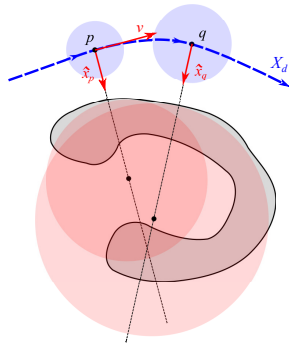
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 - ▶ Computations are parallelizable.
 - ~> *Target balls picked via edge-by-edge optimization*
- ▶ Higher-dimensional analog via *collared spheres [3]*:
 - ▶ An embedding $\gamma : \mathbb{S}^{n-1} \times \{0\} \hookrightarrow \mathbb{R}^n$ is *collared*, if it extends to an embedding of $\mathbb{S}^{n-1} \times [-1, 1] \hookrightarrow \mathbb{R}^n$.
 - ▶ *Generalized Schoenflies [4] implies*: If $\gamma : \mathbb{S}^{n-1} \hookrightarrow \mathbb{S}^n$ is collared, then $\mathbb{S}^n \setminus \gamma(\mathbb{S}^{n-1})$ is the disjoint union of two open balls.



Future Directions: Reactive TTGs

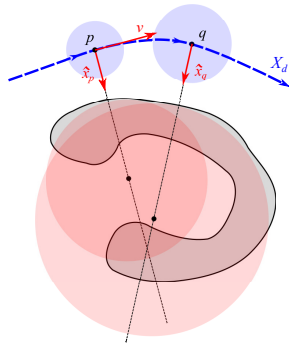
- ▶ **Returning to Single-Agent Relay Tracking,**
 - ▶ *Available plans:* For all p , $\mathcal{A}_p \equiv \mathbb{S}^{n-1}$;
 - ▶ *Desired TTG:* $f \in \mathcal{M}(\mathcal{F})$ with $f \leq \mathbf{M}_{\mathcal{F}}^u$;
 - ▶ *Task:* X_d , parametric or subdivided.





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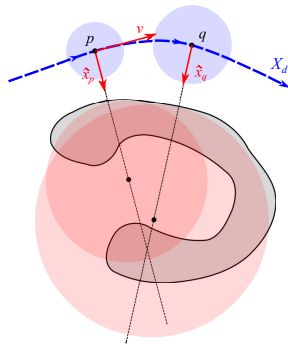
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 - ▶ f increases in the direction \hat{x}_p ;
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 - ▶ f increases in the direction \hat{x}_p ;
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- ▶ **Problem 3:** Compute approximations $f \leq \mathbf{M}_{\mathcal{F}}^U$ in closed form.
- ▶ **Problem 4:** Determine the relationship between ∇f and TTG plans.
 - ▶ Characterizing TTG plans *locally*?
 - ▶ Value tradeoffs *à-la* [5, 6] between tracking X_d and detours into \mathcal{F} ?





Future directions: Learning TTGs

In the absence of complete knowledge of \mathcal{F} ...

- ▶ Instead of learning \mathcal{F} , learn $M_{\mathcal{F}}^{\mathcal{U}}$ (or possibly $\nabla M_{\mathcal{F}}^{\mathcal{U}}$), as a model.



Future directions: Learning TTGs

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- ▶ **Instead of learning \mathcal{F} , learn $M_{\mathcal{F}}^U$ (or possibly $\nabla M_{\mathcal{F}}^U$), as a model.**
- ▶ With f_t approximating $M_{\mathcal{F}}^U$ at time t ,
 - ▶ f_t corresponding to a polygonal approximation \mathcal{F}_t of \mathcal{F} ;
 - ▶ Monotonicity, $t \leq s \implies f_t \leq f_s$, ensure a valid TTG for all time.
 - ▶ **Challenge:** Maintain a connected model of \mathcal{F} , or—
 - ▶ **Challenge:** Extend methods over disconnected \mathcal{F} , and—
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- ▶ Closed parametric form $M_{\mathcal{F}}^{\mathcal{U}} = H(c^*)$, $c^* \in \mathbb{R}^N$ means...
 - ▶ f_t takes the form of $H(c_t)$, $c_t \in \mathbb{R}^N$;
 - ▶ Think of N as a bound on the complexity of \mathcal{F}_t ;
 - ▶ **WE COULD ATTEMPT LEARNING c^* ;**
 - ▶ A natural loss function is, e.g., $\mathcal{L}(c) \triangleq \|M_{\mathcal{F}}^{\mathcal{U}} - H(c)\|_2^2$;
 - ▶ **Question:** Does GD over \mathcal{L} respond well to the above challenges?
 - ▶ **Question:** Could exploration be *guided* so that GD responds adequately to these challenges?



THANK YOU FOR YOUR ATTENTION!



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